

# BOUNDS FOR FLUID MODELS DRIVEN BY SEMI-MARKOV INPUTS

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In this paper we consider an infinite buffer fluid model whose input is driven by independent semi-Markov processes. The output capacity of the buffer is a constant. We derive upper and lower bounds for the limiting distribution of the stationary buffer content process. We discuss examples and applications where the results can be used to determine bounds on the loss probability in telecommunication networks.

## 1. INTRODUCTION

In high-speed telecommunication networks (mainly asynchronous transfer mode (ATM)), the optimal design and admission control problems frequently require com-

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puting a certain Quality of Service (QoS) or Grade of Service that the network users need in order to be assured. This QoS is based on loss probability, delay, delay-jitter, etc. We mainly focus on the loss probability aspect where loss occurs whenever a buffer overflows. It becomes important, therefore, to study the buffer content processes.

The high speed (e.g., 155–622 Mbits/sec) of the ATM network implies that it can transmit millions of cells (53-byte ATM packets) per second. This makes fluid-flow models useful in describing the flow of cells. We analyze the packetized traffic by approximating it by fluids, following the large literature using fluid-flow models for communication systems (see, e.g., Anick, Mitra, and Sondhi [1], Elwalid and Mitra [10]). Chen and Yao [5,6], Ott and Shanthikumar [25], Harrison [14], Chen and Mandelbaum [4], among others, demonstrate how to convert any discrete arrival system into a fluid-flow system and apply the fluid-model results. Our results can therefore be applied to a wide variety of networks, not just high-speed ATM networks.

The most popular method to analyze the buffer content process is by using the effective bandwidth approximation. The effective-bandwidth methodology, although simple to use, is based on an exponential approximation to the tail of the distribution of the buffer content in steady state. This approximation holds when the buffer sizes are very large and the tail probabilities are small. Several researchers have attempted to redress these shortcomings. For example, Elwalid et al. [9] and Elwalid and Mitra [11] modify the effective-bandwidth methodology and develop the Chernoff Dominant Eigenvalue (CDE) approximation for single-class traffic. To avoid approximations, other approaches have been developed. They include deriving upper and lower bounds for the tail of buffer content process in steady state with a Markov additive input by discretizing time and using extensions of Kingman's exponential bounds for waiting times in the stationary regime in a  $G/G/1$  queue (see Kingman [16], Ross [28], Artiges and Nain [2], and Liu, Nain, and Towsley [21]). Artiges and Nain [2] obtain exponential bounds for multiplexing multiclass Markovian on-off sources, where the upper bounds are similar to those in Palmowski and Rolski [26].

Liu et al. [21] obtain exponential bounds for a large class of single resource systems fed by multiplexing Markov Arrival Processes in discrete time. In this paper using the exponential change of measure, we generalize the results of Liu et al. to the continuous time case and a more general input process. Obviously, the parameters in the exponents of the lower and upper bounds obtained in both papers using a Markov additive input model are the same. The technique of exponential change of measure are presented in Palmowski and Rolski [26], who develop bounds for the distribution of the buffer-content process whose input traffic is modulated by a continuous time Markov chain (CTMC). In this paper we generalize their results to the semi-Markov modulated traffic.

The paper is organized as follows. In Section 2 we describe the model of a single buffer with traffic from  $K$  independent sources modulated by semi-Markov processes. In Section 3 we explain the notations used for semi-Markov processes for the

case where there is only a single source ( $K = 1$ ). In Section 4 we derive bounds for the buffer content distribution for the single source ( $K = 1$ ) case. In Section 5 we generalize the results in Section 4 to  $K$  sources. In Section 6 we demonstrate how to compute the bounds, and in Section 7 we illustrate the results using several examples.

**2. SINGLE BUFFER FLUID MODEL**

Consider a single buffer that admits single-class traffic from  $K$  independent sources, and the  $k$ th source driven by a random environment process  $\{Z^k(t), t \geq 0\}$ ,  $k = 1, 2, \dots, K$  (Fig. 1). Note that  $Z^k(t)$  can be thought of as the state of the  $k$ th input source at time  $t$ . When source  $k$  is in state  $Z^k(t)$ , it generates fluid at rate  $r_{Z^k(t)}^k$  into the buffer. Let  $X(t)$  be the amount of fluid in the buffer at time  $t$ . The buffer has infinite capacity and is serviced by a channel of constant capacity  $c$ . The dynamics of the buffer-content process  $\{X(t), t \geq 0\}$  are described by

$$\frac{dX(t)}{dt} = \begin{cases} \sum_{k=1}^K r_{Z^k(t)}^k - c & \text{if } X(t) > 0 \\ \left\{ \sum_{k=1}^K r_{Z^k(t)}^k - c \right\}^+ & \text{if } X(t) = 0, \end{cases} \tag{1}$$

where  $\{x\}^+ = \max(x, 0)$ . The solution is given by (Kulkarni and Rolski [20])

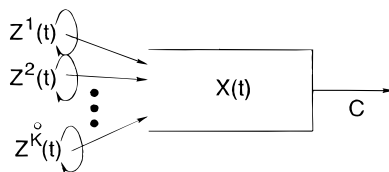
$$X(t) = \sup_{0 \leq u \leq t} \left( Y(t), \int_u^t \left( \sum_{k=1}^K r_{Z^k(s)}^k - c \right) ds \right),$$

where

$$Y(t) = X(0) + \int_0^t \left( \sum_{k=1}^K r_{Z^k(s)}^k - c \right) ds.$$

It has been shown in [20] that the buffer-content process  $\{X(t), t \geq 0\}$  is stable if

$$\sum_{k=1}^K E\{r_{Z^k(\infty)}^k\} < c, \tag{2}$$



**FIGURE 1.** Single buffer fluid model.

in which case  $X(t) \rightarrow X$  in distribution with

$$X = \sup_{u \leq 0} \int_u^0 \left( \sum_{k=1}^K r_{Z^k(s)} - c \right) ds. \tag{3}$$

One of the main aims of this paper is to analyze the limiting behavior of the buffer-content process  $\{X(t), t \geq 0\}$ , namely,

$$P\{X > x\} = \lim_{t \rightarrow \infty} P\{X(t) > x\}. \tag{4}$$

We obtain upper and lower bounds for this limiting distribution  $P\{X > x\}$ . We first consider a single source ( $K = 1$ ) and then extend the analysis to multiple sources.

### 3. SINGLE SOURCE MODEL: NOTATION

Consider the fluid model in Section 2 (see Fig. 1) with a single source ( $K = 1$ ) modulated by a semi-Markov process (SMP)  $\{Z(t), t \geq 0\}$  on state space  $\{1, 2, \dots, \ell\}$ . Fluid is generated at rate  $r_i$  at time  $t$  when the SMP is in state  $Z(t) = i$ . We first introduce the relevant notation for the SMP  $\{Z(t), t \geq 0\}$ .

Let  $S_n$  denote the time of the  $n$ th jump epoch in the SMP with  $S_0 = 0$ . Define  $Z_n$  as the state of the SMP immediately after the  $n$ th jump; that is,

$$Z_n = Z(S_n+).$$

Let

$$G_{ij}(x) = P\{S_1 \leq x; Z_1 = j | Z_0 = i\}. \tag{5}$$

The kernel of the SMP is

$$G(x) = [G_{ij}(x)]_{i,j=1,\dots,\ell}.$$

Note that  $\{Z_n, n \geq 0\}$  is a discrete time Markov chain (DTMC) with transition probability matrix

$$P = G(\infty).$$

We assume that this DTMC is irreducible and recurrent. Let

$$G_i(x) = P\{S_1 \leq x | Z_0 = i\} = \sum_{j=1}^{\ell} G_{ij}(x)$$

and the expected time the SMP spends in state  $i$  be

$$\tau_i = E(S_1 | Z_0 = i).$$

Let

$$\pi_i = \lim_{n \rightarrow \infty} P\{Z_n = i\}$$

be the stationary distribution of the DTMC  $\{Z_n, n \geq 0\}$ . It is given by the unique nonnegative solution to

$$\pi = \pi P \quad \text{and} \quad \sum_i \pi_i = 1. \tag{6}$$

The stationary distribution of the SMP is given by

$$p_i = \lim_{n \rightarrow \infty} P\{Z(t) = i\} = \frac{\pi_i \tau_i}{\sum_{m=1}^{\ell} \pi_m \tau_m}. \tag{7}$$

Since

$$E\{r_{Z(\infty)}\} = \sum_{m=1}^{\ell} p_m r_m,$$

for stability we assume that (see Eq. (2))

$$\sum_{m=1}^{\ell} p_m r_m < c. \tag{8}$$

We also assume that

$$c < \max_{1 \leq m \leq \ell} r_m; \tag{9}$$

otherwise, the buffer will always be empty in steady state.

Let  $L(t)$  be the remaining sojourn time of the SMP in the current state at time  $t$ . It is known that (see Kulkarni [17])

$$\lim_{t \rightarrow \infty} P\{Z(t) = m, L(t) > x\} = p_m \int_x^{\infty} \frac{1 - G_m(s)}{\tau_m} ds.$$

$\{L(t), t \geq 0\}$  is called the supplementary age process of  $\{Z(t), t \geq 0\}$ .

#### 4. SINGLE SOURCE MODEL: RESULTS

The following theorem gives the bounds on the limiting distribution of  $X(t)$ . First, we need the following definitions. Let

$$\phi_{ij}(\delta) = E(e^{\delta(r_i - c)S_1}; Z_1 = j | Z_0 = i), \tag{10}$$

and

$$\Phi(\delta) = [\phi_{ij}(\delta)]_{i,j \in 1, \dots, \ell}. \tag{11}$$

Note that  $\Phi(\delta)$  is a matrix with nonnegative elements. Let  $e(\Phi(\delta))$  be the Perron–Frobenius eigenvalue of  $\Phi(\delta)$ . Let  $\eta$  be the smallest real-positive value satisfying

$$e(\Phi(\eta)) = 1. \tag{12}$$

Let  $h = [h_1, h_2, \dots, h_\ell]$  be the corresponding left eigenvector, that is,

$$h = h\Phi(\eta). \tag{13}$$

Define

$$H = \sum_{i=1}^{\ell} \frac{h_i}{\eta(r_i - c)} \left( \sum_{j=1}^{\ell} (\phi_{ij}(\eta)) - 1 \right), \tag{14}$$

$$\Psi_{\max}(i, j) = \sup_x \left\{ \frac{h_i e^{-\eta(r_i - c)x} \int_x^{\infty} e^{\eta(r_i - c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} \int_x^{\infty} dG_{ij}(y)} \right\} \tag{15}$$

and

$$\Psi_{\min}(i, j) = \inf_x \left\{ \frac{h_i e^{-\eta(r_i - c)x} \int_x^{\infty} e^{\eta(r_i - c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} \int_x^{\infty} dG_{ij}(y)} \right\}. \tag{16}$$

**THEOREM 1:** *Suppose inequalities (8) and (9) hold and there exists solution of Eq. (12) (which is unique) and*

$$\Psi_{\min}(i, j) > 0 \text{ and } \Psi_{\max}(i, j) < \infty \text{ for all } i, j \in \{1, \dots, \ell\}. \tag{17}$$

*Then  $X(t) \rightarrow X$  in distribution and*

$$C_* e^{-\eta x} \leq P\{X > x\} \leq C^* e^{-\eta x}, \text{ for all } x \geq 0, \tag{18}$$

*where  $\eta$  is from Eq. (12) and*

$$C^* = \frac{H}{\min_{i:r_i > c, j:p_{ij} > 0} \{\Psi_{\min}(i, j)\}}$$

$$C_* = \frac{H}{\max_{i:r_i > c, j:p_{ij} > 0} \{\Psi_{\max}(i, j)\}}.$$

We now describe the main steps in the proof of this theorem below. Each step is discussed in detail in the following subsections.

**Step 1. The reversed SMP.** From Eq. (3) we have

$$P\{X > x\} = P\left\{ \sup_{u \leq 0} \int_u^0 (r_{Z(s)} - c) ds > x \right\}$$

$$= P\left\{ \sup_{t \geq 0} \int_0^t (r_{Z(-s)} - c) ds > x \right\}.$$

Now define

$$Y(s) = Z(-s).$$

Then  $\{Y(s), -\infty < s < \infty\}$  is the time-reversed version of  $\{Z(s), -\infty < s < \infty\}$ , therefore,

$$P\{X > x\} = P\left\{\sup_{t \geq 0} \int_0^t (r_{Y(s)} - c) ds > x\right\}.$$

Thus, we need to study the reversed process  $\{Y(t), -\infty < t < \infty\}$ . In general  $\{Y(t), -\infty < t < \infty\}$  is not an SMP. A necessary and sufficient condition for it to be an SMP is given in Section 4.1.

*Step 2. Markov process and its generator.* Even if  $\{Y(t), -\infty < t < \infty\}$  is an SMP, it is not a Markov process in general. Hence, we construct a canonical Markov process  $\{w(t), t \geq 0\}$  defined by  $w(t) = (Y(t), S(t))$ , where  $S(t)$  is the supplementary age process of  $\{Y(t), -\infty < t < \infty\}$  (as defined in Section 3). Let  $Q$  be its generator. We shall show in Section 4.2 that

$$[Q(x)]_{ij} = \begin{cases} \frac{\partial}{\partial x} - \rho_i(x) + \rho_i(x)a_{ii}\mathbf{D} & i = j, \\ \rho_i(x)a_{ij}\mathbf{D} & i \neq j. \end{cases}$$

Hence, the process defined by

$$M(t) = \frac{v(w(t))}{v(w(0))} e^{-\int_0^t \frac{Qv}{v}(w(s)) ds}$$

is a martingale (Ethier and Kurtz [12]), for any function  $v(\cdot, \cdot) \in \mathcal{D}(Q)$  (the domain of the generator  $Q$ ) with strictly positive infimum.

*Step 3. Choice of  $v(\cdot, \cdot)$ .* We next show (see Section 4.3) that it is possible to choose a  $v(\cdot, \cdot) \in \mathcal{D}(Q)$  such that

$$M(t) = \frac{v(w(t))}{v(w(0))} e^{-\int_0^t \frac{Qv}{v}(w(s)) ds} = \frac{v(w(t))}{v(w(0))} e^{\zeta \int_0^t (r_{Y(s)} - c) ds},$$

for some positive real number  $\zeta$ . Note that  $\{M(t), t \geq 0\}$  is a mean one martingale with respect to the natural filtration of  $\{w(t), t \geq 0\}$  process, that is,  $E[M(t)] = 1$  for all  $t$ .

*Step 4. Martingale  $M(t)$ .* Let

$$\tau(x) = \inf\left\{t > 0 : \int_0^t (r_{Y(s)} - c) ds > x\right\}. \tag{19}$$

From the above definition, we have,

$$\left\{ \sup_{t \geq 0} \int_0^t (r_{Y(s)} - c) ds > x \right\} \equiv \{\tau(x) < \infty\}.$$

Hence,

$$P \left\{ \sup_{t \geq 0} \int_0^t (r_{Y(s)} - c) ds > x \right\} = P\{\tau(x) < \infty\}.$$

Now define a modified measure  $\tilde{P}_t$  on  $\mathcal{F}_t = \sigma\{w(u) : 0 \leq u \leq t\}$  (see Section 4.4) by

$$\frac{d\tilde{P}_t}{dP_t} = M(t).$$

Then,

$$\begin{aligned} P\{\tau(x) < \infty\} &= \int_{\tau(x) < \infty} dP \\ &= \int_{\tau(x) < \infty} [M(\tau(x))]^{-1} d\tilde{P}. \end{aligned}$$

From Lemma 4,

$$\tilde{P}(\tau(x) < \infty) = 1. \tag{20}$$

Therefore,

$$\begin{aligned} P\{\tau(x) < \infty\} &= \int_{\tau(x) < \infty} \frac{v(w(0))}{v(w(\tau(x)))} e^{-\xi x} d\tilde{P} \\ &= e^{-\xi x} \tilde{E} \left\{ \frac{v(w(0))}{v(w(\tau(x)))}; \tau(x) < \infty \right\} \end{aligned} \tag{21}$$

$$= e^{-\xi x} \sum_{j=1}^{\ell} \int_0^{\infty} \tilde{E}_{(j,y)} \left\{ \frac{v(j,y)}{v(w(\tau(x)))} \right\} \pi_j(y) dy, \tag{22}$$

where  $\tilde{E}$  is the expectation with respect to  $\tilde{P}$  and we use Eq. (20) to justify equality between (21) and (22).

**Step 5. The bounds.** Clearly, at time  $\tau(x)$ ,  $w(\tau(x))$  can only be in a state  $(i, s)$  such that  $r_i > c$ . Hence, the lower bound on  $\tilde{E}_{(j,y)}\{[1/v(w(\tau(x)))]\}$  is  $1/\{\max_{i:r_i > c} \sup_x v(i, x)\}$  and the upper bound on  $\tilde{E}_{(j,y)}\{[1/v(w(\tau(x)))]\}$  is  $1/\{\min_{i:r_i > c} \inf_x v(i, x)\}$ . These yield the bounds (see Section 4.5 for the expressions) on  $P\{X > x\}$  in terms of the parameters of the reversed process  $Y(t)$ .

**Step 6. Original process.** We convert the bounds on  $P\{X > x\}$  in terms of that of the  $\{Z(t), -\infty < t < \infty\}$  process. This is done in Section 4.6.



**4.1. Time-Reversed Semi-Markov Processes**

As mentioned in Step 1 in Section 4, we shall use reversed SMPs in the proof of Theorem 1. Hence, we collect the relevant results about reversed SMPs here.

DEFINITION: An SMP with kernel  $G(\cdot)$  is called a nonanticipative SMP if

$$[G_{ij}(x)] = [G_i(x)p_{ij}].$$

We call such an SMP “nonanticipative” because the sojourn time in the current state does not depend upon the following state. In other words, given the current state, the sojourn time and the next state are independent of each other (a property that is exhibited by CTMCs). In case of a general SMP, the sojourn time depends on both the current state and the following state.

We first restate a theorem proved in Chari [3].

THEOREM 2: Let  $\{Z(t), -\infty < t < \infty\}$  be a semi-Markov process with kernel  $G(x) = [G_{ij}(x)]$ . Then the time-reversed process  $\{Y(t) = Z(-t), -\infty < t < \infty\}$  is an SMP if and only if  $\{Z(t), -\infty < t < \infty\}$  is a nonanticipative SMP. If  $\{Z(t), -\infty < t < \infty\}$  is a nonanticipative SMP, then the reversed SMP is also a nonanticipative SMP with kernel

$$F(x) = [F_{ij}(x)] = [F_i(x)a_{ij}], \tag{23}$$

where

$$F_i(x) = G_i(x), \tag{24}$$

$$A = [a_{ij}] = \left[ \begin{matrix} & \pi_j \\ p_{ji} & \pi_i \end{matrix} \right], \tag{25}$$

and  $\pi$  is given by Eq. (6).

Note that the stationary distribution of the reversed SMP is the same as that of the original SMP.

We will use the following notations. Let  $S_n$  denote the time of the  $n$ th jump epoch in the SMP  $\{Y(t), t \geq 0\}$ , with  $S_0 = 0$ . Define  $Y_n$  as the state of the SMP immediately after the  $n$ th jump, that is,

$$Y_n = Y(S_n+).$$

Define

$$\chi(\delta) = [\chi_{ij}(\delta)] = [E(e^{\delta(r_i-c)S_1}; Y_1 = j | Y_0 = i)]. \tag{26}$$

Note that  $\chi(\delta)$  is a matrix with nonnegative elements. Let  $e(\chi(\delta))$  be the Perron–Frobenius eigenvalue of  $\chi(\delta)$ . Let  $\zeta$  be the smallest real positive value satisfying

$$e(\chi(\zeta)) = 1. \tag{27}$$

Let

$$\Gamma(x) = [\Gamma_{ij}(x)] = [E(e^{\xi(r_i-c)(S_1-x)}; Y_1 = j | Y_0 = i, S_1 > x)]. \tag{28}$$

**4.2. Reversed Markov Process and Its Generator**

In this section we explain in detail the results shown in Step 2 of Section 4. We assume that  $\{Z(t), t \geq 0\}$  is a nonanticipative SMP. We consider its reversed (nonanticipative) semi-Markov process  $\{Y(t), t \geq 0\}$  with kernel  $F(x)$ , together with its supplementary age process  $\{S(t)\}$ . The process  $\{w(t), t \geq 0\}$  defined by

$$w(t) = (Y(t), S(t)) \tag{29}$$

is Markovian; moreover, it is a piecewise-deterministic Markov process, according to the terminology of Davis [7,8].

The generator

$$Q = [q_{ij}]_{i,j=1,\dots,\ell} \tag{30}$$

and its domain  $\mathcal{D}(Q)$  are defined as follows. We use column vector function  $f = (f_1, \dots, f_\ell)^T$ , where  $f_i: R_+ \rightarrow R$  are measurable functions, and adopt a notation that  $f(i, x) = f_i(x)$  (we use them interchangeably). Consider the Markov process  $w(t)$  on the probability space  $(\Omega, \mathcal{F}, P_{(i,y)})$ , where  $P_{(i,y)}$  is the underlying probability measure for which  $Y(0) = i, S(0) = y$ . For measurable functions  $f, f^*: R_+ \rightarrow R^\ell$  we denote

$$M_{f,f^*}(w(t)) = f_{Y(t)}(S(t)) - f_i(y) - \int_0^t f^*(Y(s), S(s)) ds, \quad t \geq 0.$$

We look for all measurable vector functions  $f, f^*$  such that  $\{M_{f,f^*}(w(t)), t \geq 0\}$  is a  $P_{(i,y)}$  martingale for all  $(i, y)$ . We denote this family of  $f$  as  $\mathcal{D}(Q)$ , and we call the mapping  $Q: f \rightarrow f^*$  as the (full) generator. The results of Theorem 26.14 from Davis [8] are adapted to the process  $w(t)$  as follows: The family  $\mathcal{D}(Q)$  consists of all measurable functions  $f(x) = (f_1(x), \dots, f_\ell(x))$ , such that  $f_i(x)$  is absolutely continuous on  $[0, s_i)$ , and  $E|f(i, T_i)| < \infty$ , where  $s_i = \inf\{t: F_i(t) = 1\}$  and  $T_i$  is a random variable with distribution function  $F_i(x)$ . Also,  $C_b(R_+)$  is the set of all continuous and bounded functions  $f: R_+ \rightarrow R$ . Then, the family of functions  $f(x) = (f_1(x), \dots, f_\ell(x)) \in C_b^\ell(R_+)$ , which are absolutely continuous on  $[0, s_i)$  are included in  $\mathcal{D}(Q)$ . As condition  $E|f(i, T_i)| < \infty$  is computationally very difficult to check, we only consider the family of functions  $C_b^\ell(R_+)$ .

Define the hazard rate function

$$\rho_i(x) = \frac{f_i(x)}{\bar{F}_i(x)},$$

where  $f_i(x) = dF_i(x)/dx$ , and  $\bar{F}_i(x) = 1 - F_i(x)$ .

Following Davis [7,8], the full generator of  $\{w(t)\}$  is as stated in the following lemma.

LEMMA 1: *The full generator of the process  $w(t)$  is  $(Qf)(i, x) = (Q(x)f(x))_i$ , where*

$$[Q(x)]_{ij} = \begin{cases} \frac{\partial}{\partial x} - \rho_i(x) + \rho_i(x)a_{ii}\mathbf{D} & i = j \\ \rho_i(x)a_{ij}\mathbf{D} & i \neq j, \end{cases} \tag{31}$$

where  $i, j = 1, \dots, \ell$  and  $(\mathbf{D}d)(x) = d(0)$  for a function  $d \in C_b(R_+)$ . Moreover,  $\mathcal{D}(Q)$  is its domain.

### 4.3. The $v$ Function

We now continue with Step 3 of Section 4 and show how to choose  $v(x) = (v_1(x), \dots, v_\ell(x)) \in \mathcal{D}(Q)$ . Following the idea from Palmowski and Rolski [27], we look for the smallest  $\beta < 0$  fulfilling

$$Q(x)v(x) = \beta\Delta v(x), \tag{32}$$

where  $\Delta = \text{diag}([r_i - c])$

$$v \in C_b^\ell(R_+) \tag{33}$$

and

$$\inf_{i,x} v_i(x) > 0. \tag{34}$$

The  $i$ th row of Eq. (32) is the first order nonhomogeneous differential equation

$$\frac{\partial}{\partial x} v_i(x) + \rho_i(x) \left( \sum_{j=1}^{\ell} a_{ij}v_j(0) - v_i(x) \right) = \beta(r_i - c)v_i(x), \tag{35}$$

whose general solution is

$$v_i(x) = (\bar{F}_i(x))^{-1} e^{\beta(r_i-c)x} \left[ v_i(0) - \int_0^x e^{-\beta(r_i-c)z} f_i(z) dz \sum_{j=1}^{\ell} a_{ij}v_j(0) \right].$$

A necessary condition for (34) is

$$v_i(0) - \int_0^\infty e^{-\beta(r_i-c)z} f_i(z) dz \sum_{j=1}^{\ell} a_{ij}v_j(0) \geq 0$$

for all  $i = 1, \dots, \ell$ . This is equivalent to  $u \geq \chi(-\beta)u$  where  $u = (v_1(0), \dots, v_\ell(0))^T$  and  $v_i(0) > 0$ .

LEMMA 2: *The smallest possible  $\beta < 0$  fulfilling*

$$Q(x)v(x) = \beta\Delta v(x)$$

is  $-\zeta$ . Then,

$$v(x) = \Gamma(x)u, \tag{36}$$

where  $\chi(\zeta)u = u$ .

PROOF: We look for the smallest  $\beta$  and vector  $u > 0$  such that  $\chi(-\beta)u \leq u$ . If there exist  $-\beta_0 > \zeta$  and  $u > 0$  fulfilling the above inequality, then multiplying it by the left Perron–Frobenius vector  $\nu_{-\beta_0} > 0$ , we would have  $e(\chi(-\beta_0))\nu_{-\beta_0}u \leq \nu_{-\beta_0}u$  so  $e(\chi(-\beta_0)) \leq 1$ . But this is a contradiction as  $e(\chi(\alpha))$  is a convex function and  $e(\chi(\zeta)) = 1$ . Thus,  $-\beta = \zeta$  and  $(v_1(0), \dots, v_\ell(0)) = (u_1, \dots, u_\ell)$ , where  $\chi(\zeta)u = u$ . In that case,

$$v_i(x) = (\bar{F}_i(x))^{-1} e^{-\zeta(r_i-c)x} \sum_{j=1}^{\ell} a_{ij}u_j \int_x^{\infty} e^{\zeta(r_i-c)z} f_i(z) dz \tag{37}$$

$$= \sum_{j=1}^{\ell} \Gamma_{ij}(x)u_j; \tag{38}$$

that is,  $v(x) = \Gamma(x)u$ . ■

#### 4.4. The Martingale $M(t)$

As mentioned in Step 4 of Section 4, we develop the mean one martingale  $M(t)$  as follows. By Proposition 3.2 in Ethier and Kurtz [12], we have that

$$M(t) = \frac{v(w(t))}{v(i, x)} e^{-\int_0^t \frac{Qv}{v}(w(s)) ds} = \frac{v(w(t))}{v(i, x)} e^{\zeta \int_0^t (r_{Y(s)}-c) ds}$$

is a martingale on  $(\Omega, \mathcal{F}, \{\mathcal{F}_t\}, P_{(i,x)})$ , where we may suppose that  $\Omega = D[0, \infty)$ ,  $\mathcal{F}_t$  is the history of  $w(t)$  up to time  $t$ , and  $\mathcal{F}$  is the smallest  $\sigma$  field generated by all subsets from  $\mathcal{F}_t$ .

Following Palmowski and Rolski [27], we define a new probability measure  $\tilde{P}_{(i,x)}$  on  $(\Omega, \mathcal{F}, \{\mathcal{F}_t^\omega\})$  by

$$\frac{d\tilde{P}_{(i,x),t}}{dP_{(i,x),t}} = M(t)$$

on which a canonical Markov process  $\{w(t)\}$  is Markovian with full generator

$$\tilde{Q}g = \frac{Q(vg)}{v} - g \frac{Qv}{v} = \frac{Q(vg)}{v} + \zeta g \Delta.$$

Thus the  $i$ th component is

$$[\tilde{Q}g](i, x) = \frac{\partial}{\partial x} g_i(x) + \frac{\rho_i(x)}{v_i(x)} \sum_{j=1}^{\ell} a_{ij}u_j \left[ \sum_{j=1}^{\ell} \frac{a_{ij}u_j}{\sum_k a_{ik}u_k} g_j(0) - g_i(x) \right]. \tag{39}$$

LEMMA 3: *On the new probability space  $(\Omega, \mathcal{F}, \{\mathcal{F}_t\}, \tilde{P}_{(i,x)})$ , the process  $\{Y(t), -\infty < t < \infty\}$  is again a semi-Markov process specified by  $([\tilde{a}_{ij}], [\tilde{\rho}_i(x)], [r_i])$ , where*

$$\tilde{a}_{ij} = \frac{a_{ij}u_j}{\sum_{k=1}^{\ell} a_{ik}u_k} = \frac{\chi_{ij}(\zeta)u_j}{u_i}, \tag{40}$$

$$\tilde{\rho}_i(x) = \frac{\rho_i(x) \sum_{j=1}^{\ell} a_{ij}u_j}{v_i(x)} = \frac{f_i(x)e^{\zeta(r_i-c)x} \sum_{j=1}^{\ell} a_{ij}u_j}{\sum_{j=1}^{\ell} \int_x^{\infty} e^{\zeta(r_i-c)z} a_{ij}f_i(z)u_j dz}, \tag{41}$$

$$\tilde{f}_i(x) = \frac{f_i(x)e^{\zeta(r_i-c)x} \sum_{j=1}^{\ell} a_{ij}u_j}{u_i}, \tag{42}$$

and

$$\tilde{\tau}_i = \frac{(\chi'(\zeta)u)_i}{u_i(r_i - c)}. \tag{43}$$

PROOF: The first parts of Eqs. (40) and (41) are simple consequences of the definition of generator  $Q$  and the function  $v_i(x)$ . Moreover, the second part of Eq. (40) follows from

$$\begin{aligned} \frac{a_{ij}u_j}{\sum_{k=1}^{\ell} a_{ik}u_k} &= \frac{a_{ij} \int_0^{\infty} e^{\zeta(r_i-c)x} f_i(x) dx u_j}{\sum_{i=1}^{\ell} a_{ik} \int_0^{\infty} e^{\zeta(r_i-c)x} f_i(x) dx u_k} \\ &= \frac{\chi_{ij}(\zeta)u_j}{\sum_{k=1}^{\ell} \chi(\zeta)_{ik}u_k} \\ &= \frac{\chi_{ij}(\zeta)u_j}{u_i}. \end{aligned}$$

From Eq. (41), we get

$$\tilde{f}_i(x) = \frac{f_i(x)e^{\zeta(r_i-c)x} \sum_{j=1}^{\ell} a_{ij}u_j}{d_i}, \tag{44}$$

where  $d_i$  is a normalizing constant. Hence,

$$d_i = \int_0^\infty f_i(z) e^{\zeta(r_i - c)z} \sum_{j=1}^\ell a_{ij} u_j dz = \sum_{j=1}^\ell \chi(\zeta)_{ij} u_j = u_i \tag{45}$$

and we obtain Eq. (42). Using Eq. (42), Eq. (43) is fulfilled by

$$\begin{aligned} (r_i - c) \tilde{\tau}_i &= (r_i - c) \int_0^\infty x \tilde{f}_i(x) dx \\ &= \frac{\sum_{k=1}^\ell \chi'_{ik}(\zeta) u_k}{u_i} \\ &= \frac{(\chi'(\zeta) u)_i}{u_i}. \end{aligned}$$

This completes the proof of this lemma. ■

Let  $\tilde{\pi} \tilde{P} = \tilde{\pi}$  be the stationary distribution of the Markov chain  $\{Y_n\}$  under the new probability measure. In matrix notation, Eq. (40) reads as

$$\tilde{P} = [\tilde{a}_{ij}] = \text{diag}(u^{-1}) \chi(\zeta) \text{diag}(u). \tag{46}$$

Therefore, from Eq. (46)

$$\tilde{\pi} = \tilde{\pi} \text{diag}(u^{-1}) \chi(\zeta) \text{diag}(u);$$

hence,

$$\tilde{\pi} = \nu \text{diag}(u), \tag{47}$$

where  $\nu$  is the left (row) eigenvector of  $\chi(\zeta)$  corresponding to eigenvalue 1 such that  $\nu u = 1$ ; that is,

$$\nu = \nu \chi(\zeta).$$

The drift  $\tilde{d} = \tilde{E}Y(t) - c$  is

$$\tilde{d} = \sum_{i=1}^\ell \frac{r_i \tilde{\pi}_i \tilde{\tau}_i}{\sum_{j=1}^\ell \tilde{\pi}_j \tilde{\tau}_j} - c.$$

Hence,  $\tilde{d} > 0$  if and only if

$$\sum_{i=1}^\ell \tilde{\pi}_i \tilde{\tau}_i (r_i - c) > 0,$$

which is, by Eqs. (43) and (47), equivalent to

$$\nu \chi'(\zeta) u > 0. \tag{48}$$

Consider the eigenvalue problem

$$\begin{cases} \chi(\alpha) h_\alpha = \kappa(\alpha) h_\alpha \\ \nu_\alpha \chi(\alpha) = \kappa(\alpha) \nu_\alpha \\ \nu_\alpha h_\alpha = 1 \\ \alpha \geq 0. \end{cases} \tag{49}$$

Function  $\kappa(\alpha)$  is convex (by Kingman–Miller theorem; see Miller [22] and the beautiful proof of this theorem by Kingman [15]).

LEMMA 4:  $\tilde{d} > 0$ .

PROOF: Differentiating the first line in Eq. (49) we obtain

$$\chi'(\alpha) h_\alpha + \chi(\alpha) h'_\alpha = \kappa'(\alpha) h_\alpha + \kappa(\alpha) h'_\alpha.$$

Multiplying both the sides from the right by  $\nu_\alpha$  and rearranging we arrive at

$$\nu_\alpha \chi'(\alpha) h_\alpha = \kappa'(\alpha) + (\kappa(\alpha) - 1) \nu_\alpha h'_\alpha. \tag{50}$$

Since  $\chi(0) = P$ , we have

$$h_0 = Pe = e, \quad \nu_0 = \pi P = \pi, \quad \kappa(0) = 1, \tag{51}$$

where  $e = [1, \dots, 1]^T$  is the (column) vector of ones. The stability condition  $d = \sum_i p_i r_i - c < 0$  and Eq. (50) yield

$$\begin{aligned} \kappa'(0) &= \pi \chi'(0) e \\ &= \sum_{i=1}^{\ell} a_i \sum_{j=1}^{\ell} \frac{\partial}{\partial \alpha} \chi_{ij}(\alpha(r_i - c)) \Big|_{\alpha=0} \\ &= \sum_{i=1}^{\ell} a_i \tau_i (r_i - c) < 0. \end{aligned}$$

Because  $\kappa(0) = 1$ ,  $\kappa'(0) < 0$ , and  $\kappa(\alpha)$  is a convex function, there is  $\kappa'(\zeta) > 0$  for  $\kappa(\zeta) = 1$ . Substituting  $\alpha = \zeta$  in Eq. (50) and bearing in mind inequality (48),  $\tilde{d} > 0$  is equivalent to

$$\nu \chi'(\zeta) u = \kappa'(\zeta) > 0.$$

The proof is completed. ■

4.5. The Bounds for the Reversed SMP

Continuing with the analysis in Step 5 of Section 4, consider a reversed nonanticipative SMP  $\{Y(t), \infty < t < \infty\}$  with kernel  $F(x) = [F_{ij}(x)] = [F_i(x)a_{ij}]$ . Define

$$\pi_m(x) = a_m \frac{1 - F_m(x)}{\tau_m},$$

where  $a_m$  is the stationary probability of the reversed SMP  $\{Y(t), -\infty < t < \infty\}$ . Let

$$b_m = \frac{a_m}{\tau_m \zeta (r_m - c)}. \tag{52}$$

**THEOREM 3:** *If conditions (33) and (34) are fulfilled, then the steady-state distribution of the buffer-content process is bounded as*

$$C_* e^{-\zeta x} \leq P\{X > x\} \leq C^* e^{-\zeta x}, \tag{53}$$

where

$$C^* = \frac{b(I - A)u}{\min_{m:r_m > c} \inf_{x \geq 0} v_m(x)}$$

$$C_* = \frac{b(I - A)u}{\max_{m:r_m > c} \sup_{x \geq 0} v_m(x)}.$$

**PROOF:** Let  $\tau(x)$  be as in Eq. (19). Continuing from Eq. (22), we have

$$P\{\tau(x) < \infty\} = e^{-\zeta x} \sum_{j=1}^{\ell} \int_0^{\infty} \tilde{E}_{(j,y)} \left\{ \frac{v_j(y)}{v(w(\tau(x)))} \right\} \pi_j(y) dy,$$

which can be bounded above by

$$\leq e^{-\zeta x} \frac{1}{\min_{m:r_m > c} \inf_{x \geq 0} v_m(x)} \sum_{j=1}^{\ell} \int_0^{\infty} v_j(y) \pi_j(y) dy. \tag{54}$$

However,

$$\begin{aligned} \sum_{j=1}^{\ell} \int_0^{\infty} v_j(y) \pi_j(y) dy &= \sum_{j=1}^{\ell} \frac{a_j}{\tau_j} \int_0^{\infty} v_j(y) \bar{F}_j(y) dy \\ &= \sum_{j=1}^{\ell} \frac{a_j}{\tau_j} \int_0^{\infty} \sum_{k=1}^{\ell} a_{jk} u_k \int_y^{\infty} e^{\zeta(r_j - c)z} f_j(z) dz e^{-\zeta(r_j - c)y} dy, \end{aligned}$$



and, integrating by parts,

$$\begin{aligned}
 &= \sum_{j=1}^{\ell} \left[ \frac{a_j}{\tau_j \zeta(r_j - c)} \sum_{k=1}^{\ell} a_{jk} u_k [\tilde{F}_j(-\zeta(r_j - c)) - 1] \right] \\
 &= \sum_{j=1}^{\ell} b_j \sum_{k=1}^{\ell} a_{jk} u_k \tilde{F}_j(-\zeta(r_j - c)) - \sum_{j=1}^{\ell} b_j \sum_{k=1}^{\ell} a_{jk} u_k \\
 &= \sum_{j=1}^{\ell} b_j (\chi(\zeta)u)_j - \sum_{j=1}^{\ell} b_j (Au)_j = b(I - A)u,
 \end{aligned}$$

where  $\tilde{F}(\alpha)$  is a Laplace transform. Thus, by Eq. (54), the proof of the upper bound in Eq. (53) is completed. In a similar way we obtain the lower bounds in Eq. (53). ■

**4.6. Bounds for Nonanticipative SMP**

In the previous subsection we derived bounds for the buffer-content process in terms of the parameters of the reversed process  $\{Y(t), -\infty < t < \infty\}$ , as illustrated in Theorem 3. In this subsection, following Step 6 of Section 4, we convert the bounds in Theorem 3 in terms of the original (nonanticipative) SMP  $\{Z(t), -\infty < t < \infty\}$ . We assume that  $\{Z(t), -\infty < t < \infty\}$  is an  $\ell$ -state nonanticipative semi-Markov process with kernel  $[G_{ij}(x)] = [G_i(x)p_{ij}]$ . All other parameters of the SMP (viz.  $\pi_j, p_j, \tau_j, \phi_{ij}(\delta), \eta, h_j$ ) are as described in Sections 3 and 4.

From Eqs. (25) and (24) in Theorem 2, we have

$$a_{ij} = p_{ji} \frac{\pi_j}{\pi_i} \tag{55}$$

and

$$F_i(x) = G_i(x).$$

From the definition of  $\chi_{ij}(\delta)$  in Eq. (26) and Eq. (10), we obtain

$$\begin{aligned}
 \chi_{ij}(\delta) &= \tilde{F}_{ij}(-\delta(r_i - c)) \\
 &= a_{ij} \tilde{F}_i(-\delta(r_i - c)) \\
 &= \tilde{G}_i(-\delta(r_i - c)) p_{ji} \frac{\pi_j}{\pi_i},
 \end{aligned} \tag{56}$$

and

$$\begin{aligned}
 \phi_{ij}(\delta) &= \tilde{G}_{ij}(-\delta(r_i - c)) \\
 &= \tilde{G}_i(-\delta(r_i - c)) p_{ij}.
 \end{aligned} \tag{57}$$

LEMMA 5: For a given real positive number  $\delta$ , the matrices  $\chi(\delta)$  and  $\Phi(\delta)$  have identical eigenvalues. Therefore,

$$e(\Phi(\delta)) = e(\chi(\delta))$$

and  $\eta = \zeta$ .

PROOF: Let  $\mathcal{G}$  and  $\Delta$  be diagonal matrices such that

$$\mathcal{G} = \text{diag}\{\tilde{G}_1(-\delta(r_1 - c)), \tilde{G}_2(-\delta(r_2 - c)), \dots, \tilde{G}_\ell(-\delta(r_\ell - c))\}$$

and

$$\Delta = \text{diag}\{\pi_1, \pi_2, \dots, \pi_\ell\}.$$

Clearly,

$$\Phi(\delta) = \mathcal{G}P$$

and

$$\chi(\delta) = \mathcal{G}\Delta^{-1}P'\Delta.$$

Let  $\lambda$  be an eigenvalue of  $\Phi(\delta)$ . Therefore, if  $|D|$  denotes the determinant of the square matrix  $D$  and if  $D'$  denotes the transpose of a matrix  $D$ , we have

$$\begin{aligned} |\mathcal{G}P - \lambda I| &= 0, \\ \Rightarrow |\mathcal{G}||P - \lambda\mathcal{G}^{-1}| &= 0, \\ \Rightarrow |\mathcal{G}||P' - \lambda\mathcal{G}^{-1}| &= 0, \\ \Rightarrow |\mathcal{G}||\Delta^{-1}||\Delta||\Delta^{-1}P'\Delta - \lambda\mathcal{G}^{-1}| &= 0, \\ \Rightarrow |\mathcal{G}\Delta^{-1}P'\Delta - \lambda I| &= 0. \end{aligned}$$

Thus,  $\lambda$  is also an eigenvalue of  $\chi(\delta)$ . For a given real positive number  $\delta$ , the matrices  $\chi(\delta)$  and  $\Phi(\delta)$  have identical eigenvalues. Specifically, the largest real positive eigenvalues of  $\chi(\delta)$  and  $\Phi(\delta)$  are identical. Also, the smallest  $\delta$  that satisfies the largest real positive eigenvalues of  $\chi(\delta)$  and  $\Phi(\delta)$  to be identical and equal to one is unique. Therefore,

$$e(\Phi(\delta)) = e(\chi(\delta))$$

and  $\eta = \zeta$ . ■

Since  $\eta = \zeta$ , the eigenvectors  $h$  and  $u$  are given by

$$h\Phi(\eta) = h \quad \text{and} \quad \chi(\eta)u = u.$$

LEMMA 6: The eigenvector  $u$ , the vector  $v(x)$ , and the matrix  $\Gamma(x)$  are related to the corresponding terms of the original nonanticipative SMP  $\{Z(t), -\infty < t < \infty\}$  as given below:

$$u_j = \frac{\tilde{G}_j(-\eta(r_j - c))}{\pi_j} h_j, \tag{58}$$

$$\Gamma_{ij}(x) = \frac{\int_x^\infty e^{\eta(r_i - c)y} dG_i(y)}{\int_x^\infty e^{\eta(r_i - c)x} dG_i(y)} p_{ji} \frac{\pi_j}{\pi_i}, \tag{59}$$

and

$$v_i(x) = \frac{\int_x^\infty e^{\eta(r_i - c)y} dG_i(y)}{\int_x^\infty e^{\eta(r_i - c)x} dG_i(y)} \frac{h_i}{\pi_i}. \tag{60}$$

PROOF: Since  $\chi(\eta)u = u$ , from Eq. (56) for all  $i$ , we get

$$\begin{aligned} u_i &= \sum_{j=1}^\ell \chi_{ij}(\eta) u_j \\ &= \sum_{j=1}^\ell \tilde{G}_i(-\eta(r_i - c)) p_{ji} \frac{\pi_j}{\pi_i} u_j. \end{aligned} \tag{61}$$

Note that

$$u_j = \frac{\tilde{G}_j(-\eta(r_j - c))}{\pi_j} h_j$$

satisfies Eq. (61) since

$$\begin{aligned} \sum_{j=1}^\ell \tilde{G}_i(-\eta(r_i - c)) p_{ji} \frac{\pi_j}{\pi_i} u_j &= \frac{\tilde{G}_i(-\eta(r_i - c))}{\pi_i} \sum_{j=1}^\ell p_{ji} \pi_j u_j \\ &= \frac{\tilde{G}_i(-\eta(r_i - c))}{\pi_i} \sum_{j=1}^\ell \tilde{G}_j(-\eta(r_j - c)) p_{ji} h_j \\ &= \frac{\tilde{G}_i(-\eta(r_i - c))}{\pi_i} \sum_{j=1}^\ell \phi_{ji}(\eta) h_j \\ &= \frac{\tilde{G}_i(-\delta(r_i - c))}{\pi_i} h_i \\ &= u_i \end{aligned}$$

using Eq. (57) and the relation  $h\Phi(\eta) = h$ . From Eq. (28),

$$\begin{aligned} \Gamma_{ij}(x) &= \frac{\int_x^\infty e^{\zeta(r_i-c)y} dF_i(y)}{\int_x^\infty e^{\zeta(r_i-c)x} dF_i(y)} a_{ij} \\ &= \frac{\int_x^\infty e^{\eta(r_i-c)y} dG_i(y)}{\int_x^\infty e^{\eta(r_i-c)x} dG_i(y)} p_{ji} \frac{\pi_j}{\pi_i}. \end{aligned} \tag{62}$$

Note that from Eq. (36),

$$v(x) = \Gamma(x)u. \tag{63}$$

Using Eq. (62) and

$$h_i = \sum_{j=1}^{\ell} \tilde{G}_j(-\eta(r_j - c)) p_{ji} h_j,$$

we have

$$\begin{aligned} v_i(x) &= \sum_{j=1}^{\ell} \frac{\int_x^\infty e^{\eta(r_i-c)y} dG_i(y)}{\int_x^\infty e^{\eta(r_i-c)x} dG_i(y)} p_{ji} \tilde{G}_j(-\eta(r_j - c)) \frac{h_j}{\pi_i} \\ &= \frac{\int_x^\infty e^{\eta(r_i-c)y} dG_i(y)}{\int_x^\infty e^{\eta(r_i-c)x} dG_i(y)} \frac{h_i}{\pi_i}. \end{aligned} \quad \blacksquare$$

From the definition of  $b_i$  in Eq. (52), it follows that

$$b_i = \frac{p_i}{\tau_i \eta(r_i - c)}. \tag{64}$$

The following theorem describes the bounds of Theorem 3 in terms of the original nonanticipative SMP  $\{Z(t), -\infty < t < \infty\}$ .

THEOREM 4: *The steady-state distribution of the buffer-content process is bounded as*

$$C_* e^{-\eta x} \leq P(X > x) \leq C^* e^{-\eta x},$$

where

$$C^* = \frac{\sum_{i=1}^{\ell} \frac{h_i}{\eta(r_i - c)} (\tilde{G}_i(-\eta(r_i - c)) - 1)}{\min_{i:r_i > c} \inf_x \left\{ \frac{h_i \tau_i}{p_i} \frac{\int_x^{\infty} e^{\eta(r_i - c)y} dG_i(y)}{\int_x^{\infty} e^{\eta(r_i - c)x} dG_i(y)} \right\}} \tag{65}$$

and

$$C_* = \frac{\sum_{i=1}^{\ell} \frac{h_i}{\eta(r_i - c)} (\tilde{G}_i(-\eta(r_i - c)) - 1)}{\max_{i:r_i > c} \sup_x \left\{ \frac{h_i \tau_i}{p_i} \frac{\int_x^{\infty} e^{\eta(r_i - c)y} dG_i(y)}{\int_x^{\infty} e^{\eta(r_i - c)x} dG_i(y)} \right\}} \tag{66}$$

provided that

$$C_* > 0 \quad \text{and} \quad C^* < \infty. \tag{67}$$

PROOF: Using the time-reversed process  $\{Y(t), -\infty < t < \infty\}$ , we get from Theorem 3

$$C^* = \frac{b(I - A)u}{\min_{i:r_i > c} \inf_x v_i(x)} \tag{68}$$

$$C_* = \frac{b(I - A)u}{\max_{i:r_i > c} \sup_x v_i(x)}. \tag{69}$$

Now we rewrite  $b, A, u,$  and  $v_i(x)$  in Eqs. (68) and (69) in terms of the parameters of the  $\{Z(t), -\infty < t < \infty\}$  process using Eqs. (64), (55), (58), and (60). After some algebra, we obtain Eqs. (65) and (66). Let us note that conditions (33) and (34) are equivalent assumptions that

$$C_* > 0 \quad \text{and} \quad C^* < \infty. \tag{70}$$

Hence, the proof is complete. ■

**4.7. General SMP Proof**

Theorem 4 of the previous section holds only for the special case when the SMP driving the input to a buffer can be modeled as a nonanticipative SMP. There are several applications where the nonanticipative SMP models fail. In this section, we prove Theorem 1 (which is stated for a general SMP) using Theorem 4 for the nonanticipative SMP. We first state and prove a theorem that explains how to convert a general SMP into a nonanticipative SMP.

Let  $\{Z(t), t \geq 0\}$  be an  $\ell$ -state (not necessarily nonanticipative) SMP with kernel  $[G_{ij}(x)]$ , transition probabilities  $p_{ij} = G_{ij}(\infty)$ , and stationary probabilities  $\pi_j$  (from  $\pi = \pi G(\infty)$ ). Also,  $S_n$  and  $Y_n$  denote respectively the  $n$ th transition epoch and the state of the SMP immediately after the  $n$ th transition; that is,  $Y_n = Z(S_n+)$ . Let  $N(t)$  be the number of transitions by the SMP until time  $t$ . Define

$$\bar{Z}(t) = (Z(S_{N(t)}+), Z(S_{N(t)+1}+)). \tag{71}$$

Let  $\bar{S}_n$  be the  $n$ th transition epoch of the  $\{\bar{Z}(t), t \geq 0\}$  process and  $\bar{Y}_n$  be the state of the  $\{\bar{Z}(t), t \geq 0\}$  process immediately after the  $n$ th transition, that is,

$$\bar{Y}_n = \bar{Z}(\bar{S}_n).$$

Observe that  $\bar{S}_n = S_n$  and  $\bar{Y}_n = (Y_n, Y_{n+1})$ .

**THEOREM 5:** *The  $\ell^2$ -state process  $\{\bar{Z}(t), t \geq 0\}$  is a nonanticipative SMP with state space  $\{(i, j) : 1 \leq i, j \leq \ell\}$ , kernel*

$$\bar{G}_{(i,j),(k,l)}(x) = \begin{cases} p_{kl} G_{ij}(x)/G_{ij}(\infty) & \text{if } j = k \\ 0 & \text{if } j \neq k, \end{cases} \tag{72}$$

*transition probabilities*

$$\bar{p}_{(i,j),(k,l)} = \begin{cases} p_{kl} & \text{if } j = k \\ 0 & \text{if } j \neq k, \end{cases} \tag{73}$$

*and stationary probabilities*

$$\bar{\pi}_{(i,j)} = \pi_i p_{ij}. \tag{74}$$

**PROOF:** Let

$$\begin{aligned} \bar{G}_{(i,j),(k,l)}(x) &= P\{\bar{Y}_{n+1} = (k, l), \bar{S}_{n+1} \leq x | \bar{Y}_n = (i, j), \bar{S}_n, \bar{Y}_{n-1}, \dots\} \\ &= P\{Y_{n+2} = l, Y_{n+1} = k, S_{n+1} \\ &\quad \leq x | Y_{n+1} = j, Y_n = i, S_n, Y_{n-1}, \dots\} \\ &= P\{Y_2 = l, Y_1 = k, S_1 \leq x | Y_1 = j, Y_0 = i\}, \end{aligned}$$

since  $\{Z(t), t \geq 0\}$  is an SMP. Also, given  $Y_1, Y_2$  is conditionally independent of  $S_1$ . Hence, we can write

$$\begin{aligned} \bar{G}_{(i,j),(k,l)}(x) &= P\{Y_2 = l, Y_1 = k | Y_1 = j, Y_0 = i\} P\{S_1 \leq x | Y_1 = j, Y_0 = i\} \\ &= \begin{cases} p_{kl} \cdot P\{S_1 \leq x | Y_1 = j, Y_0 = i\} & \text{if } j = k \\ 0 \cdot P\{S_1 \leq x | Y_1 = j, Y_0 = i\} & \text{if } j \neq k \end{cases} \\ &= \bar{p}_{(i,j),(k,l)} \frac{P\{S_1 \leq x, Y_1 = j | Y_0 = i\}}{P\{Y_1 = j | Y_0 = i\}} \\ &= \bar{p}_{(i,j),(k,l)} \frac{G_{ij}(x)}{G_{ij}(\infty)} \\ &= \bar{p}_{(i,j),(k,l)} \bar{G}_{(i,j)}(x), \end{aligned}$$

which is of the form (like the nonanticipative SMP)

$$[G_{ab}(x)] = [G_a(x)p_{ab}].$$

Hence, we get Eqs. (72) and (73). It is easy to verify that  $\bar{\pi}_{(i,j)} = \pi_i p_{ij}$  satisfies

$$\bar{\pi} \bar{G}(\infty) = \bar{\pi} \quad \text{and} \quad \sum_{i=1}^{\ell} \sum_{j=1}^{\ell} \bar{\pi}_{(i,j)} = 1.$$

Hence, we get Eq. (74). ■

Let  $\{Z(t), t \geq 0\}$  be a general (not necessarily nonanticipative) SMP. Construct a nonanticipative SMP  $\{\bar{Z}(t), t \geq 0\}$  as described in Eq. (71). Now we use Theorem 4 for the  $\ell^2$ -state nonanticipative SMP  $\{\bar{Z}(t), t \geq 0\}$ . Similar to the Eqs. (11) and (5), we now define  $\bar{\Phi}(v)$ ,  $\bar{r}_{(i,j)}$ , and  $\bar{G}_{(i,j)}(\cdot)$ .

From the definition of the nonanticipative SMP  $\{\bar{Z}(t), t \geq 0\}$  and Eqs. (73), (72), and (74), it is straightforward to show that

$$\bar{r}_{(i,j)} = \begin{cases} r_i & \text{if } p_{ij} > 0 \\ 0 & \text{if } p_{ij} = 0, \end{cases} \tag{75}$$

$$\begin{aligned} \bar{G}_{(i,j)}(v(\bar{r}_{(i,j)} - c)) &= \frac{\bar{G}_{ij}(v(r_i - c))}{G_{ij}(\infty)} \\ &= \frac{\phi_{ij}(v)}{p_{ij}}. \end{aligned} \tag{76}$$

We have,

$$\begin{aligned} \bar{\phi}_{(i,j),(k,l)}(v) &= E(e^{v(\bar{r}_{(i,j)} - c)\bar{S}_1}; \bar{Y}_1 = (k, l) | \bar{Y}_0 = (i, j)) \\ &= \begin{cases} 0 & \text{if } k \neq j \\ \bar{G}_{(i,j)}(-v(\bar{r}_{(i,j)} - c)) \bar{p}_{(i,j),(j,l)} & \text{if } k = j \end{cases} \\ &= \begin{cases} 0 & \text{if } k \neq j \\ \frac{\phi_{ij}(v)}{p_{ij}} p_{kl} & \text{if } k = j. \end{cases} \end{aligned} \tag{77}$$

LEMMA 7: *The smallest real positive value of  $\delta$  that satisfies*

$$e(\Phi(\delta)) = 1$$

*is  $\eta$  if and only if  $\eta$  is the smallest real positive value that satisfies*

$$e(\bar{\Phi}(\delta)) = 1.$$

PROOF: For a given  $v$ , let  $\lambda_i$  be such that

$$w^i \Phi(v) = \lambda_i w^i,$$

where  $\lambda_i$  is the  $i$ th eigenvalue of  $\Phi(v)$  and  $w^i$  is the corresponding eigenvector. Therefore  $\lambda_1, \lambda_2, \dots, \lambda_\ell$  are the  $\ell$  eigenvalues of  $\Phi(v)$  and  $w^1, w^2, \dots, w^\ell$  the corresponding eigenvectors. Without loss of generality we can assume that

$$\text{Re}(\lambda_1) \geq \text{Re}(\lambda_2) \geq \dots \geq \text{Re}(\lambda_\ell),$$

where  $\text{Re}(x)$  denotes the real part of the complex number  $x$ .

From the expression of  $\bar{\Phi}(v)$  in Eq. (77), it follows that there are at most  $\ell$  linearly independent rows for the matrix  $\bar{\Phi}(v)$ . Hence, at least  $\ell^2 - \ell$  eigenvalues are zero. Let  $\bar{\lambda}_1, \bar{\lambda}_2, \dots, \bar{\lambda}_\ell$  be the  $\ell$  possibly nonzero eigenvalues of  $\bar{\Phi}(v)$  and  $\bar{w}^1, \bar{w}^2, \dots, \bar{w}^\ell$  the corresponding eigenvectors. Without loss of generality we can assume that

$$\text{Re}(\bar{\lambda}_1) \geq \text{Re}(\bar{\lambda}_2) \geq \dots \geq \text{Re}(\bar{\lambda}_\ell).$$

From Perron–Frobenius theorem,  $\bar{\lambda}_1$  is real and positive, hence  $\text{Re}(\bar{\lambda}_1) \geq 0$ . Note that we can choose each eigenvector  $\bar{w}^n$  such that

$$\bar{w}_{(i,j)}^n = p_{ij} w_i^n.$$

Then also for a given  $v$ ,  $\bar{\lambda}_n = \lambda_n$  for all  $n \in 1, 2, \dots, \ell$ . In fact,

$$\begin{aligned} [\bar{w}^n \bar{\Phi}(v)]_{(k,l)} &= \sum_{j=1}^{\ell} \sum_{i=1}^{\ell} \bar{w}_{(i,j)}^n \bar{\phi}_{(i,j),(k,l)}(v) \\ &= \sum_{i=1}^{\ell} \bar{w}_{(i,k)}^n \bar{\phi}_{(i,k),(k,l)}(v) \\ &= \sum_{i=1}^{\ell} \bar{w}_{(i,k)}^n \frac{\phi_{ik}(v)}{p_{ik}} p_{kl} \\ &= \sum_{i=1}^{\ell} p_{ik} w_i^n \frac{\phi_{ik}(v)}{p_{ik}} p_{kl} \\ &= \lambda_n p_{kl} w_k^n \\ &= \lambda_n \bar{w}_{(k,l)}^n \\ &= \bar{\lambda}_n \bar{w}_{(k,l)}^n. \end{aligned}$$



When  $v = \eta$ , the largest real-positive eigenvalue for both  $\bar{\Phi}(\eta)$  and  $\Phi(\eta)$  are equal to one, hence, the proof. ■

In the next lemma, we derive the relationship between the left eigenvectors  $h$  and  $\bar{h}$ .

LEMMA 8: Let  $\bar{h}\bar{\Phi}(\eta) = \bar{h}$  and  $h\Phi(\eta) = h$ , then

$$\bar{h}_{(i,j)} = h_i p_{ij}. \tag{78}$$

PROOF: From the definition of  $\bar{h}$ , we have,

$$\begin{aligned} \bar{h}_{(k,l)} &= \sum_{j=1}^{\ell} \sum_{i=1}^{\ell} \bar{h}_{(i,j)} \bar{\phi}_{(i,j),(k,l)}(\eta) \\ &= \sum_{i=1}^{\ell} \bar{h}_{(i,k)} \bar{\phi}_{(i,k),(k,l)}(\eta) \\ &= \sum_{i=1}^{\ell} \bar{h}_{(i,k)} \frac{\phi_{ik}(\eta)}{p_{ik}} p_{kl}. \end{aligned}$$

Since  $h_j = \sum_{i=1}^{\ell} \phi_{ij}(\eta) h_i$ , verify that

$$\bar{h}_{(i,j)} = h_i p_{ij} \tag{79}$$

solves  $\bar{h}\bar{\Phi}(\eta) = \bar{h}$ . ■

Define  $\bar{\tau}_{(i,j)}$  as the expected time the nonanticipative SMP  $\{\bar{Z}(t), t \geq 0\}$  spends in state  $(i, j)$ . Therefore,

$$\begin{aligned} \bar{\tau}_{(i,j)} &= E\{\bar{S}_1 | \bar{Z}_0 = (i, j)\} \\ &= \int_0^{\infty} (1 - \bar{G}_{(i,j)}(x)) dx \\ &= \int_0^{\infty} (1 - G_{ij}(x)/G_{ij}(\infty)) dx. \end{aligned}$$

Let  $\bar{p}_{(i,j)}$  be the probability that, in steady-state, the nonanticipative SMP  $\{\bar{Z}(t), t \geq 0\}$  is state  $(i, j)$  and is given by

$$\bar{p}_{(i,j)} = \frac{\bar{\tau}_{(i,j)} \bar{\pi}_{(i,j)}}{\sum_{k=1}^{\ell} \sum_{l=1}^{\ell} \bar{\tau}_{(k,l)} \bar{\pi}_{(k,l)}}.$$

Subsequently, we need an expression for  $\bar{\tau}_{(i,j)}p_{ij}/\bar{p}_{(i,j)}$ . This is shown in the following lemma.

LEMMA 9:

$$\frac{\bar{\tau}_{(i,j)}p_{ij}}{\bar{p}_{(i,j)}} = \frac{\tau_i}{p_i}. \tag{80}$$

PROOF:

$$\begin{aligned} \frac{\bar{\tau}_{(i,j)}p_{ij}}{\bar{p}_{(i,j)}} &= \frac{\bar{\tau}_{(i,j)}\pi_i p_{ij}}{\bar{p}_{(i,j)}\pi_i} \\ &= \frac{\bar{\tau}_{(i,j)}\bar{\pi}_{(i,j)}}{\bar{p}_{(i,j)}\pi_i} \\ &= \frac{1}{\pi_i} \sum_{k=1}^{\ell} \sum_{l=1}^{\ell} \bar{\tau}_{(k,l)}\bar{\pi}_{(k,l)} \\ &= \frac{1}{\pi_i} \sum_{k=1}^{\ell} \sum_{l=1}^{\ell} \pi_k p_{kl} \int_0^{\infty} (1 - G_{kl}(x)/G_{kl}(\infty)) dx \\ &= \frac{1}{\pi_i} \sum_{k=1}^{\ell} \pi_k \int_0^{\infty} (1 - G_k(x)) dx \\ &= \frac{1}{\pi_i} \sum_{k=1}^{\ell} \pi_k \tau_k \\ &= \frac{\tau_i}{p_i}. \end{aligned} \quad \blacksquare$$

Next, we give the proof of Theorem 1.

PROOF OF THEOREM 1: From Eq. (65) for the nonanticipative SMP  $\{\bar{Z}(t), t \geq 0\}$ , we have

$$\bar{C}^* = \frac{\sum_{i=1}^{\ell} \sum_{j=1}^{\ell} \frac{\bar{h}_{(i,j)}}{\eta(\bar{r}_{(i,j)} - c)} (\bar{G}_{(i,j)}(-\eta(\bar{r}_{(i,j)} - c)) - 1)}{\min_{i,j:\bar{r}_{(i,j)} > c} \inf_x \left\{ \frac{\bar{h}_{(i,j)}\bar{\tau}_{(i,j)}}{\bar{p}_{(i,j)}} \frac{\int_x^{\infty} e^{\eta(\bar{r}_{(i,j)} - c)y} d\bar{G}_{(i,j)}(y)}{\int_x^{\infty} e^{\eta(\bar{r}_{(i,j)} - c)x} d\bar{G}_{(i,j)}(y)} \right\}}. \tag{81}$$

We substitute the expressions  $\bar{h}_{(i,j)}$ ,  $\bar{r}_{(i,j)}$ ,  $\bar{G}_{(i,j)}(\cdot)$ ,  $\bar{G}_{(i,j)}(y)$ , and  $\bar{\tau}_{(i,j)}p_{ij}/\bar{p}_{(i,j)}$  from Eqs. (79), (75), (76), (72), and (80), respectively, into Eq. (81) to get

$$\begin{aligned} \bar{C}^* &= \frac{\sum_{i=1}^{\ell} \sum_{j=1}^{\ell} \frac{h_i p_{ij}}{\eta(r_i - c)} (\bar{G}_{ij}(-\eta(r_i - c))/p_{ij} - 1)}{\min_{i,j:r_i > c, p_{ij} > 0} \inf_x \left\{ \frac{h_i \tau_i \int_x^{\infty} e^{\eta(r_i - c)y} dG_{ij}(y)}{p_i \int_x^{\infty} e^{\eta(r_i - c)x} dG_{ij}(y)} \right\}} \\ &= \frac{\sum_{i=1}^{\ell} \frac{h_i}{\eta(r_i - c)} \left( \sum_{j=1}^{\ell} (\phi_{ij}(\eta)) - 1 \right)}{\min_{i,j:r_i > c, p_{ij} > 0} \inf_x \left\{ \frac{h_i e^{-\eta(r_i - c)x} \int_x^{\infty} e^{\eta(r_i - c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} \int_x^{\infty} dG_{ij}(y)} \right\}} \\ &= C^*. \end{aligned}$$

Note that assumption (17) yields (67). Using a similar analysis for  $\bar{C}_*$ , Theorem 1 is proved. ■

**5. MULTIPLEXING OF INDEPENDENT SOURCES**

In this section we consider the model in Section 2. As illustrated in Figure 1, there are  $K$  independent sources, each modulated by an SMP  $\{Z^k(t), t \geq 0\}$  on state space  $\{1, 2, \dots, \ell_k\}$ . Fluid is generated from source  $k$  at rate  $r_i^k$  at time  $t$  when its modulating SMP is in state  $i$ . The kernel of the SMP modulating the  $k$ th input source is  $G^k(x) = [G_{ij}^k(x)]$ . The expected time the  $k$ th SMP spends in state  $i$  is  $\tau_i^k$ . The stationary distributions of the  $k$ th SMP  $\{Z^k(t), t \geq 0\}$  and its underlying DTMC are, respectively,  $p^k$  and  $\pi^k$ . We assume for stability and nontriviality that

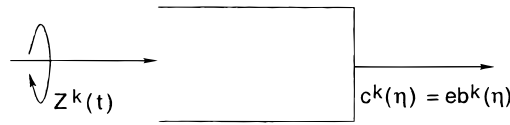
$$\sum_{k=1}^K \sum_{m=1}^{\ell_k} p_m^k r_m^k < c < \sup_{k,m} r_m^k.$$

Let  $A_k(t)$  be the total amount of fluid input from the  $k$ th source to the buffer over time  $(0, t]$ , where

$$A_k(t) = \int_0^t r_{Z^k(u)}^k du. \tag{82}$$

Then the effective bandwidth of the  $k$ th source,  $eb^k(\cdot)$ , is given by

$$eb^k(v) = \lim_{t \rightarrow \infty} \frac{1}{vt} \log E\{\exp(vA_k(t))\}. \tag{83}$$



**FIGURE 2.** Single source modification.

Let  $\eta$  be the smallest positive solution to

$$\sum_{k=1}^K eb^k(\eta) = c. \tag{84}$$

As a first step we analyze a single source model where source  $k$  acts as the only input source to an infinite-capacity buffer with output capacity

$$c^k(\eta) = eb^k(\eta) \tag{85}$$

(see Fig. 2). For this single source model, following the analysis in Section 4, for each source  $k$ , we compute  $\Phi^k(\eta)$  and  $h^k$  from Eqs. (10) and (13) as

$$\phi_{ij}^k(\eta) = E(e^{\eta(r_i^k - c^k(\eta))S_1^k}; Z_1^k = j | Z_0^k = i), \tag{86}$$

$$h^k = h^k\Phi^k(\eta). \tag{87}$$

Define

$$P^k(i, j) = P\{Z_1^k = j | Z_0^k = i\}. \tag{88}$$

Continuing with the single source model, we also define

$$H^k = \sum_{i=1}^{\ell_k} \frac{h_i^k}{\eta(r_i^k - c^k(\eta))} \left( \sum_{j=1}^{\ell_k} (\phi_{ij}^k(\eta)) - 1 \right), \tag{89}$$

$$\Psi_{\min}^k(i, j) = \inf_x \left\{ \frac{h_i^k e^{-\eta(r_i^k - c^k(\eta))x} \int_x^\infty e^{\eta(r_i^k - c^k(\eta))y} dG_{ij}^k(y)}{\frac{p_i^k}{\tau_i^k} \int_x^\infty dG_{ij}^k(y)} \right\}, \tag{90}$$

and

$$\Psi_{\max}^k(i, j) = \sup_x \left\{ \frac{h_i^k e^{-\eta(r_i^k - c^k(\eta))x} \int_x^\infty e^{\eta(r_i^k - c^k(\eta))y} dG_{ij}^k(y)}{\frac{p_i^k}{\tau_i^k} \int_x^\infty dG_{ij}^k(y)} \right\}. \tag{91}$$

Using these quantities, the bounds on the multiplexed traffic are given in the following theorem.

**THEOREM 6:** *If  $\Psi_{\min}^k(i, j) > 0$  and  $\Psi_{\max}^k(i, j) < \infty$  for all  $i, j \in \{1, \dots, \ell_k\}$ ,  $k \in \{1, \dots, K\}$ , then the bounds on the limiting distribution of the buffer-content process  $\{X(t), t \geq 0\}$  driven by  $K$  independent SMP sources are given by*

$$C_* e^{-\eta x} \leq P(X > x) \leq C^* e^{-\eta x},$$

where  $\eta$  is given by the solution to Eq. (84),

$$C^* = \frac{\prod_{k=1}^K H^k}{\min_{\mathcal{A}} \prod_{k=1}^K \Psi_{\min}^k(i_k, j_k)},$$

$$C_* = \frac{\prod_{k=1}^K H^k}{\max_{\mathcal{A}} \prod_{k=1}^K \Psi_{\max}^k(i_k, j_k)},$$

and

$$\mathcal{A} = \left\{ (i_1, j_1), (i_2, j_2), \dots, (i_K, j_K) : \sum_{k=1}^K r_{i_k}^k > c \quad \text{and} \quad \forall k, \quad P^k(i_k, j_k) > 0 \right\}.$$

(92)

**PROOF:** We describe only a brief outline of the proof of the theorem by illustrating the main steps of the proof, as most of the analysis is identical to the proof of Theorem 1.

*Step 1. The nonanticipative SMPs.* For each  $k \in [1, 2, \dots, K]$ , we first construct an  $\ell_k^2$  state nonanticipative SMP  $\{\bar{Z}^k(t), t \geq 0\}$  from the SMP  $\{Z^k(t), t \geq 0\}$  following the technique in Section 4.7.

*Step 2. The reversed SMP.* Next we construct the reversed SMP of  $\{\bar{Z}^k(t), t \geq 0\}$  as

$$Y^k(s) = \bar{Z}^k(-s).$$

Therefore,

$$P\{X > x\} = P\left\{ \sup_{t \geq 0} \int_0^t \left( \sum_{k=1}^K r_{Y^k(s)}^k - c \right) ds > x \right\}.$$

*Step 3. Markov process and its generator.* We construct a Markov process  $\{W(t), t \geq 0\}$  defined by

$$W(t) = (Y^1(t), S^1(t), Y^2(t), S^2(t), \dots, Y^K(t), S^K(t)),$$

where  $S^k(t)$  is the supplementary age process of  $\{Y^k(t), -\infty < t < \infty\}$  (as defined in Section 3). Let  $Q^k$  ( $1 \leq k \leq K$ ) be the full generator of  $\{w^k(t)\} = \{Y^k(t), S^k(t)\}$  as in Eq. (29).

*Step 4. Martingale  $M(t)$ .* Let  $\eta$  satisfy Eq. (84) and for this  $\eta$  let  $v^k(i, x)$  satisfy

$$v^k(i, x) = [\Gamma^k(x)u^k]_i,$$

where

$$\Gamma_{ij}^k(x) = E(e^{\eta(r_i^k - c^k(\eta))(S_1^k - x)}; Y_1^k = j | Y_0^k = i, S_1^k > x)$$

(see Eq. (28)) and  $u^k$  is the right eigenvector satisfying

$$\chi^k(\eta)u^k = u^k$$

with  $\chi_{ij}^k(\eta) = E(e^{\eta(r_i^k - c^k(\eta))S_1^k}; Y_1^k = j | Y_0^k = i)$  (see Eq. (26)). Assume that

$$v^k \in C_b(R_+) \quad \text{and} \quad \inf_x v^k(i, x) > 0. \tag{93}$$

Then we have  $(Q^k v^k)(i, x) = -\eta(r_i^k - c)v^k(i, x)$  and, following Proposition 3.2 from Ethier and Kurtz [12], we define exponential martingales by

$$M^k(t) = \frac{v^k(w(t))}{v^k(i, x)} e^{-\eta \int_0^t (Y^k(s) - c^k(\eta)) ds}, \quad k = 1, \dots, K.$$

Since sources are independent, the process

$$\begin{aligned} M(t) &= \prod_{k=1}^K M^k(t) = \frac{v^1(w^1(t)) \cdots v^K(w^K(t))}{v^1(i^1, x^1) \cdots v^K(i^K, x^K)} e^{-\eta \int_0^t \sum_{k=1}^K (Y^k(s) - c^k(\eta)) ds} \\ &= \frac{v(W(t))}{v(i, x)} e^{-\eta \int_0^t \sum_{k=1}^K (Y^k(s) - c^k(\eta)) ds} \\ &= \frac{v(W(t))}{v(i, x)} e^{-\eta \int_0^t \left( \sum_{k=1}^K Y^k(s) - c \right) ds} \end{aligned}$$

is a martingale, where  $v(i, x) = v(i^1, x^1, \dots, i^K, x^K) = \prod_{k=1}^K v^k(i^k, x^k)$  and  $i = (i^1, \dots, i^K)$ ,  $i^k \in \{1, \dots, \ell_k\}$  and  $x = (x^1, \dots, x^K) \in R_+^K$ .

Then we also construct a new probability space  $(\Omega, \mathcal{F}, \{\mathcal{F}_t\}, \tilde{P})$  by

$$\frac{d\tilde{P}_t}{dP_t} = M(t). \tag{94}$$

*Step 5. First passage time  $\tau(x)$ .* Let

$$\tau(x) = \inf \left\{ t > 0 : \int_0^t \left( \sum_{k=1}^K r_{Y^k(s)}^k - c \right) ds > x \right\}. \tag{95}$$

From the above definition, we have,

$$\left\{ \sup_{t \geq 0} \int_0^t \left( \sum_{k=1}^K r_{Y^k(s)}^k - c \right) ds > x \right\} \equiv \{ \tau(x) < \infty \}.$$

Hence,

$$P \left\{ \sup_{t \geq 0} \int_0^t \left( \sum_{k=1}^K r_{Y^k(s)}^k - c \right) ds > x \right\} = P \{ \tau(x) < \infty \}.$$

Then, by Eq. (94),

$$\begin{aligned} P \{ \tau(x) < \infty \} &= \int_{\tau(x) < \infty} dP \\ &= \int_{\tau(x) < \infty} [M(\tau(x))]^{-1} d\tilde{P}. \end{aligned}$$

Lemma 4 yields that  $\tilde{E} r_{Z^k(\infty)}^k > c^k(\eta)$ . Hence,  $\tilde{P}(\tau(x) < \infty) = 1$  and then

$$\begin{aligned} P \{ \tau(x) < \infty \} &= \int_{\tau(x) < \infty} \frac{v(W(0))}{v(W(\tau(x)))} e^{-\eta x} d\tilde{P} \\ &= e^{-\eta x} \tilde{E} \left\{ \frac{v(W(0))}{v(W(\tau(x)))} \right\}, \end{aligned} \tag{96}$$

where  $\tilde{E}$  is the expectation with respect to  $\tilde{P}$ . Define

$$\begin{aligned} \pi_j^k &= \lim_{n \rightarrow \infty} P \{ Y_n^k = j \} \\ A_{ij}^k &= P \{ Y_1^k = j | P \{ Y_0^k = i \} \\ p_i^k &= \lim_{t \rightarrow \infty} P \{ Y^k(t) = i \} \\ \tau_i^k &= E \{ S_1^k | Y_0^k = i \} \\ b_i^k &= \frac{p_i^k}{\tau_i^k \eta (r_i^k - c^k)}. \end{aligned}$$

Next, we bound

$$\begin{aligned} \tilde{E} \left\{ \frac{v(W(0))}{v(W(\tau(x)))} \right\} &= v(W(0)) \tilde{E} \left\{ \frac{1}{v(W(\tau(x)))} \right\} \\ &= \prod_{k=1}^K \sum_{j_k=1}^{\ell_k^2} \int_0^\infty v^k(j^k, y^k) \pi_{j_k}^k(y_k) dy_k \tilde{E} \left\{ \frac{1}{v(W(\tau(x)))} \right\} \\ &= \prod_{k=1}^K b^k (I - A^k) u^k \tilde{E} \left\{ \frac{1}{v(W(\tau(x)))} \right\}. \end{aligned}$$

At time  $\tau(x)$ , the  $W(t)$  process can only be in states  $(i_1, s_1, i_2, s_2, \dots, i_K, s_K)$  that satisfy  $\sum_{k=1}^K r_{i_k}^k > c$ . Hence, we have

$$\tilde{E} \left\{ \frac{1}{v(W(\tau(x)))} \right\} \leq \frac{1}{\min_{i_1, i_2, \dots, i_K: \sum_{k=1}^K r_{i_k}^k > c} \prod_{k=1}^K \left\{ \inf_x v_k(i^k, x) \right\}}$$

and

$$\tilde{E} \left\{ \frac{1}{v(W(\tau(x)))} \right\} \geq \frac{1}{\max_{i_1, i_2, \dots, i_K: \sum_{k=1}^K r_{i_k}^k > c} \prod_{k=1}^K \left\{ \sup_x v_k(i^k, x) \right\}}.$$

Therefore the bounds on Eq. (96) are

$$L_* e^{-\eta x} \leq P\{\tau(x) < \infty\} \leq L^* e^{-\eta x}, \tag{97}$$

where

$$L^* = \frac{\prod_{k=1}^K b^k (I - A^k) u^k}{\min_{i_1, i_2, \dots, i_K: \sum_{k=1}^K r_{i_k}^k > c} \prod_{k=1}^K \left\{ \inf_x v_k(i^k, x) \right\}}$$

$$L_* = \frac{\prod_{k=1}^K b^k (I - A^k) u^k}{\max_{i_1, i_2, \dots, i_K: \sum_{k=1}^K r_{i_k}^k > c} \prod_{k=1}^K \left\{ \sup_x v_k(i^k, x) \right\}}.$$

*Step 6. Original process.* Next we rewrite  $L^*$  and  $L_*$  in terms of the  $\ell_k^2$  state SMPs  $\{\bar{Z}^k(t), t \geq 0\}$  (for all  $k \in \{1, 2, \dots, K\}$ ). Finally, we convert the expressions to those of the original  $\ell_k$  state SMPs  $\{Z^k(t), t \geq 0\}$  and obtain  $C^*$ ,  $C_*$ , and Eq. (92) of Theorem 6. Note that condition (93) put on the functions  $v^k(i, x)$  is equivalent to the condition that  $L_* > 0$  and  $L^* < \infty$ , that is, that  $C_* > 0$  and  $C^* < \infty$ , which follows from the assumptions of Theorem 6. ■

### 6. COMPUTATION OF $C^*$ AND $C_*$

In this section we address two concerns while using Theorem 1 in Section 4:

- how to compute  $\eta$  in Eq. (12), hence  $h$  in Eq. (13)
- how to compute  $\Psi_{\max}(i, j)$  and  $\Psi_{\min}(i, j)$  in Eqs. (15) and (16), respectively.



**6.1. Computation of  $\eta$  and  $h$**

To compute  $\eta$ , we have to solve Eq. (12). In terms of the effective bandwidth of the source, it is the same as solving for  $\eta$  in

$$eb(\eta) = c \tag{98}$$

whenever  $e(\Phi(v)) = 1$  has a solution (see Gautam [13]). Also, condition (17) is satisfied whenever  $e(\Phi(v)) = 1$  has a solution. Therefore, the easiest way to calculate  $\eta$  is to use Eq. (98). Since a semi-Markov process is a special case of a Markov regenerative process (MRGP), to compute the effective bandwidth of a source modulated by an SMP we use the results from Kulkarni [18].

The eigenvector  $h$  is obtained by solving

$$h\Phi(\eta) = h.$$

Note: An interesting scenario that arises in practical situations is the case when  $e(\Phi(v)) = 1$  has no solutions. This is also the case when condition (17) is not satisfied. We do not address this issue in this paper; however, we will publish a technique to get around this scenario in a forthcoming paper.

**6.2. Computation of  $\Psi_{\max}$  and  $\Psi_{\min}$**

Consider a nonnegative random variable  $Y$  with distribution  $G_{ij}(x)/G_{ij}(\infty)$  and density

$$g_{ij}(x) = \frac{dG_{ij}(x)}{dx} \frac{1}{G_{ij}(\infty)}.$$

The failure rate function of  $Y$  is defined by

$$\lambda_{ij}(x) = \frac{g_{ij}(x)}{1 - \frac{G_{ij}(x)}{G_{ij}(\infty)}}. \tag{99}$$

$Y$  is said to be an increasing failure rate (IFR) random variable if

$$\lambda_{ij}(x) \uparrow x$$

and  $Y$  is said to be a decreasing failure rate (DFR) random variable if

$$\lambda_{ij}(x) \downarrow x.$$

It is possible to obtain closed form algebraic expressions for  $\Psi_{\max}(i, j)$  and  $\Psi_{\min}(i, j)$  in Eqs. (15) and (16), respectively, if a random variable  $Y$  with distribution  $G_{ij}(x)/G_{ij}(\infty)$  is an IFR or DFR random variable. The following theorem describes how to compute  $\Psi_{\max}(i, j)$  and  $\Psi_{\min}(i, j)$  in those cases. Let  $x^*$  and  $x_*$  be such that

$$x^* = \arg \sup_x \left\{ \frac{h_i \int_x^\infty e^{\eta(r_i-c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} e^{\eta(r_i-c)x} \int_x^\infty dG_{ij}(y)} \right\} \tag{100}$$

and

$$x_* = \arg \inf_x \left\{ \frac{h_i \int_x^\infty e^{\eta(r_i-c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} e^{\eta(r_i-c)x} \int_x^\infty dG_{ij}(y)} \right\}. \tag{101}$$

**THEOREM 7:** *If Y is IFR or DFR, then  $\Psi_{\max}(i, j)$  and  $\Psi_{\min}(i, j)$  in Eqs. (15) and (16), respectively, occur at x values given by the following table:*

	IFR		DFR	
	$r_i > c$	$r_i \leq c$	$r_i > c$	$r_i \leq c$
$x^*$	0	$\infty$	$\infty$	0
$\Psi_{\max}(i, j)$	$\frac{\phi_{ij}(-\eta(r_i - c))\tau_i h_i}{p_{ij}p_i}$	$\frac{\tau_i h_i \lambda_{ij}(\infty)}{p_i(\lambda_{ij}(\infty) - \eta(r_i - c))}$	$\frac{\tau_i h_i \lambda_{ij}(\infty)}{p_i(\lambda_{ij}(\infty) - \eta(r_i - c))}$	$\frac{\phi_{ij}(-\eta(r_i - c))\tau_i h_i}{p_{ij}p_i}$
$x_*$	$\infty$	0	0	$\infty$
$\Psi_{\min}(i, j)$	$\frac{\tau_i h_i \lambda_{ij}(\infty)}{p_i(\lambda_{ij}(\infty) - \eta(r_i - c))}$	$\frac{\phi_{ij}(-\eta(r_i - c))\tau_i h_i}{p_{ij}p_i}$	$\frac{\phi_{ij}(-\eta(r_i - c))\tau_i h_i}{p_{ij}p_i}$	$\frac{\tau_i h_i \lambda_{ij}(\infty)}{p_i(\lambda_{ij}(\infty) - \eta(r_i - c))}$

where

$$\lambda_{ij}(\infty) = \lim_{x \rightarrow \infty} \lambda_{ij}(x).$$

**PROOF:** Let  $Y^t$  be the remaining life associated with the random variable  $Y$ . Then  $Y^t$  will have distribution  $G_{ij}^t(a)$  given by

$$\begin{aligned} 1 - G_{ij}^t(a) &= P\{Y^t > a\} \\ &= P\{Y - t > a | Y > t\}. \end{aligned}$$

It is known (Ross [29]) that if  $Y$  is IFR (DFR), then  $Y^t$  is stochastically decreasing (increasing) in  $t$ , hence  $E[f(Y^t)]$  is decreasing (increasing) in  $t$  for all increasing functions  $f$ . Therefore, if  $Y$  is IFR (DFR) and  $r_i > c$ , then

$$E(e^{\eta(r_i-c)[Y-x]} | Y > x) = \frac{e^{-\eta(r_i-c)x} \int_x^\infty e^{\eta(r_i-c)y} dG_{ij}(y)}{\int_x^\infty dG_{ij}(y)}$$

is decreasing (increasing) in  $x$ . Similarly, if  $Y$  is IFR (DFR), then  $Y'$  is stochastically decreasing (increasing) in  $t$ , hence  $E[f(Y')]$  is increasing (decreasing) in  $t$  for all decreasing functions  $f$ . Therefore, if  $Y$  is IFR (DFR) and  $r_i \leq c$ , then

$$E(e^{\eta(r_i-c)[Y-x]} | Y > x) = \frac{e^{-\eta(r_i-c)x} \int_x^\infty e^{\eta(r_i-c)y} dG_{ij}(y)}{\int_x^\infty dG_{ij}(y)}$$

is increasing (decreasing) in  $x$ . Note that

$$\lim_{x \rightarrow 0} \left\{ \frac{h_i \int_x^\infty e^{\eta(r_i-c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} e^{\eta(r_i-c)x} \int_x^\infty dG_{ij}(y)} \right\} = \frac{\tilde{\Phi}_{ij}(-\eta(r_i-c)) \tau_i h_i}{p_{ij} p_i}$$

and using L'Hospital's Rule we show that

$$\begin{aligned} & \lim_{x \rightarrow \infty} \left\{ \frac{h_i \int_x^\infty e^{\eta(r_i-c)y} dG_{ij}(y)}{\frac{p_i}{\tau_i} e^{\eta(r_i-c)x} \int_x^\infty dG_{ij}(y)} \right\} \\ &= \lim_{x \rightarrow \infty} \left\{ \frac{-e^{\eta(r_i-c)x} \frac{dG_{ij}(x)}{dx} \frac{h_i \tau_i}{p_i}}{-e^{\eta(r_i-c)x} \frac{dG_{ij}(x)}{dx} + \eta(r_i-c) e^{\eta(r_i-c)x} (G_{ij}(\infty) - G_{ij}(x))} \right\} \\ &= \frac{h_i \tau_i}{p_i} \lim_{x \rightarrow \infty} \left\{ \frac{p_{ij} g_{ij}(x)}{p_{ij} g_{ij}(x) - \eta(r_i-c)(p_{ij} - G_{ij}(x))} \right\} \\ &= \frac{\tau_i h_i \lambda_{ij}(\infty)}{p_i (\lambda_{ij}(\infty) - \eta(r_i-c))}. \end{aligned}$$

Hence, the result follows. ■

If  $Y$  is not an IFR or DFR random variable, then  $x^*, x_*, \Psi_{\max}(i, j)$ , and  $\Psi_{\min}(i, j)$  can be obtained only numerically.

## 7. EXAMPLES

### 7.1. General On-Off Source

Consider a source modulated by a two-state (on and off) process that alternates between on and off states. The random amount of time the process spends in the on state (called *on-times*) has cdf  $U(\cdot)$  with mean  $\tau_U$  and the corresponding *off-time* cdf

is  $D(\cdot)$  with mean  $\tau_D$ . The successive on- and off-times are independent and on-times are independent of off-times. Fluid is generated continuously at rate  $r$  during the on state and at rate 0 during the off state. A source modulated by such a two-state (on-off) process is called a *general on-off source with on-time distribution  $U(\cdot)$ , off-time distribution  $D(\cdot)$ , and (peak) rate  $r$* .

Consider a general on-off source with on-time distribution  $U(\cdot)$ , off-time distribution  $D(\cdot)$ , and rate  $r$  that inputs traffic into an infinite-capacity buffer. The output capacity of the buffer is a constant  $c$ . Assume

$$\frac{r\tau_U}{\tau_U + \tau_D} < c < r.$$

Equation (11) reduces to

$$\Phi(v) = \begin{bmatrix} 0 & \tilde{D}(vc) \\ \tilde{U}(-v(r-c)) & 0 \end{bmatrix},$$

where  $\tilde{U}(\cdot)$  and  $\tilde{D}(\cdot)$  are the Laplace Stieltjes transforms (LSTs) of  $U(t)$  and  $D(t)$ , respectively. In this subsection, we assume that  $e(\Phi(\eta)) = 1$  has a solution and it implies that

$$e(\Phi(\eta)) = \sqrt{\tilde{U}(-\eta(r-c))\tilde{D}(\eta c)} = 1$$

can be solved. Hence,  $\eta$  is the smallest real-positive solution to

$$\tilde{U}(-\eta(r-c))\tilde{D}(\eta c) = 1.$$

Also, Eq. (13) reduces to

$$h = [1 \quad \tilde{D}(\eta c)].$$

Furthermore, from Eqs. (14), (16), and (15), we get

$$H = \frac{(1 - \tilde{D}(\eta c))r}{c(r - c)},$$

$$\Psi_{\min} = \begin{bmatrix} 0 & \inf_x \left\{ \frac{(\tau_U + \tau_D) \int_x^\infty e^{-\eta c(y-x)} dD(y)}{1 - D(x)} \right\} \\ \inf_x \left\{ \tilde{D}(\eta c) \frac{(\tau_U + \tau_D) \int_x^\infty e^{\eta(r-c)(y-x)} dU(y)}{1 - U(x)} \right\} & 0 \end{bmatrix}$$

and

$$\Psi_{\max} = \begin{bmatrix} 0 & \sup_x \left\{ \frac{(\tau_U + \tau_D) \int_x^\infty e^{-\eta c(y-x)} dD(y)}{1 - D(x)} \right\} \\ \sup_x \left\{ \tilde{D}(\eta c) \frac{(\tau_U + \tau_D) \int_x^\infty e^{\eta(r-c)(y-x)} dU(y)}{1 - U(x)} \right\} & 0 \end{bmatrix}.$$

From Theorem 1, we can derive that

$$C_* e^{-\eta x} \leq P\{X > x\} \leq C^* e^{-\eta x}, \tag{102}$$

where

$$C^* = \frac{\tilde{U}(-\eta(r-c)) - 1}{\tau_U + \tau_D} \frac{r}{c(r-c)\eta \inf_x \left\{ \frac{\int_x^\infty e^{\eta(r-c)(y-x)} dU(y)}{1-U(x)} \right\}} \tag{103}$$

and

$$C_* = \frac{\tilde{U}(-\eta(r-c)) - 1}{\tau_U + \tau_D} \frac{r}{c(r-c)\eta \sup_x \left\{ \frac{\int_x^\infty e^{\eta(r-c)(y-x)} dU(y)}{1-U(x)} \right\}}. \tag{104}$$

If  $U(\cdot)$  is IFR/DFR, the supremums and the infimums in the above equations can be obtained from Theorem 7 in Section 6.2.

### 7.2. Erlang On-Off Source

Consider the special case of the general on-off source in Section 7.1 with *Erlang*  $(N_U, \alpha)$  on-time distribution, *Erlang*  $(N_D, \beta)$  off-time distribution, and rate  $r$ . Note that  $\tau_U = N_U/\alpha$  and  $\tau_D = N_D/\beta$ . Equation (11) reduces to

$$\Phi(v) = \begin{bmatrix} 0 & \tilde{D}(vc) \\ \tilde{U}(-v(r-c)) & 0 \end{bmatrix} = \begin{bmatrix} 0 & \left(\frac{\beta}{\beta+vc}\right)^{N_D} \\ \left(\frac{\alpha}{\alpha-v(r-c)}\right)^{N_U} & 0 \end{bmatrix}.$$

It is possible to show that  $e(\Phi(v)) = 1$  always has a solution. Hence,  $\eta$  is obtained by solving

$$e(\Phi(\eta)) = \sqrt{\tilde{U}(-\eta(r-c))\tilde{D}(\eta c)} = 1$$

and

$$h = \left[ 1 \left( \frac{\beta}{\beta + \eta c} \right)^{N_D} \right].$$

Using the fact that the *Erlang* random variable has an increasing hazard rate function, from Theorem 7 in Section 6.2, we see that

$$\Psi_{\min} = \begin{bmatrix} 0 & (N_U/\alpha + N_D/\beta) \left( \frac{\beta}{\beta + \eta c} \right)^{N_D} \\ \left( \frac{\beta}{\beta + \eta c} \right)^{N_D} (N_U/\alpha + N_D/\beta) \frac{\alpha}{\alpha - \eta(r-c)} & 0 \end{bmatrix}$$

and

$$\Psi_{\max} = \begin{bmatrix} 0 & (N_U/\alpha + N_D/\beta) \frac{\beta}{\beta + \eta c} \\ \left(\frac{\beta}{\beta + \eta c}\right)^{N_D} (N_U/\alpha + N_D/\beta) \left(\frac{\alpha}{\alpha - \eta(r - c)}\right)^{N_U} & 0 \end{bmatrix}.$$

Using Theorem 1, we have

$$C_* e^{-\eta x} \leq P\{X > x\} \leq C^* e^{-\eta x}, \tag{105}$$

where

$$C^* = \frac{\left(\frac{\alpha}{\alpha - \eta(r - c)}\right)^{N_U} - 1}{\tau_U + \tau_D} \frac{r}{c(r - c)\eta \left\{ \frac{\alpha}{\alpha - \eta(r - c)} \right\}} \tag{106}$$

and

$$C_* = \frac{\left(\frac{\alpha}{\alpha - \eta(r - c)}\right)^{N_U} - 1}{\tau_U + \tau_D} \frac{r}{c(r - c)\eta \left\{ \left(\frac{\alpha}{\alpha - \eta(r - c)}\right)^{N_U} \right\}}. \tag{107}$$

Consider a numerical example of an Erlang on-off source with on-time distribution  $Erlang(N_U, \alpha)$  and off-time distribution  $Erlang(N_D, \beta)$ . Let  $r = 15$ ,  $c = 10$ ,  $\tau_U = \frac{1}{70}$ , and  $\tau_D = \frac{1}{30}$ . We keep the means constant (i.e.,  $\tau_U$  and  $\tau_D$  are held constant) but decrease the variances by increasing  $N_U$  and  $N_D$ . In Figure 3 we illustrate for four pairs of  $(N_U, N_D)$ , (namely, (1,1), (4,3), (9,8), and (16,14)), the logarithm of the upper and lower bounds on the limiting distribution of the buffer-content process.

From the figure we notice that, as the variance decreases, the bounds move further apart. Also note that  $C^*$  increases with decrease in variance and  $C_*$  decreases with decrease in variance. Since  $\eta$  increases with decrease in variance, the tail of the limiting distribution rapidly goes to zero.

*Remark:* Let  $N_U \rightarrow \infty$ ,  $\alpha \rightarrow \infty$  such that  $N_U/\alpha = \tau_U$ , a finite positive number and  $N_D \rightarrow \infty$ ,  $\beta \rightarrow \infty$  such that  $N_D/\beta = \tau_D$ , a finite positive number. Such a source is a deterministic on-off source with on-times  $\tau_U$  and off-times  $\tau_D$ . The upper and lower bounds for this limiting case of an Erlang on-off source is illustrated in Figure 4.

In Figure 5, we demonstrate the exact probabilities  $P(X > x)$  for different  $c$  values increasing from  $r\tau_U/(\tau_U + \tau_D)$  to  $r$ . We compare these exact probabilities with the bounds obtained by limiting case of the Erlang distribution. The  $x^*$  values in Figures 4 and 5 are identical; hence, we can conclude that the limiting case of the Erlang distribution does produce bounds that make sense for the deterministic on-off source.

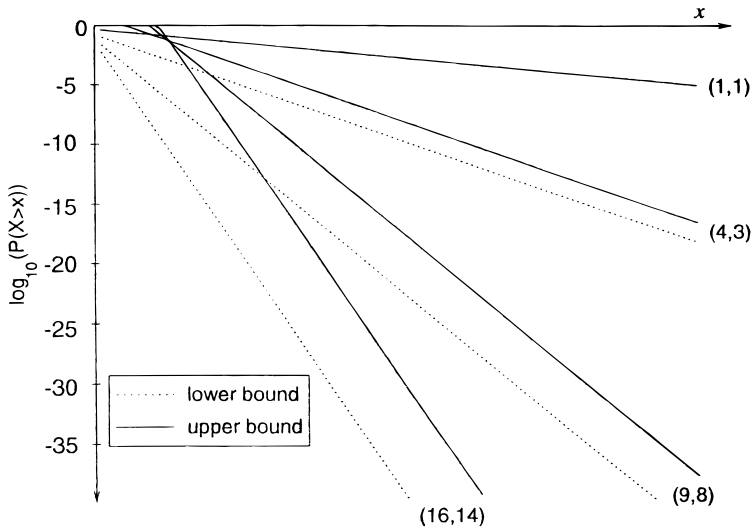


FIGURE 3. Logarithm of the upper and lower bounds as a function of  $x$ .

### 7.3. Tandem Buffers—Single Source

In this section, an exponential on-off source with on-time parameter  $\alpha$ , off-time parameter  $\beta$ , and rate  $r$  inputs traffic to an infinite-capacity buffer with output ca-

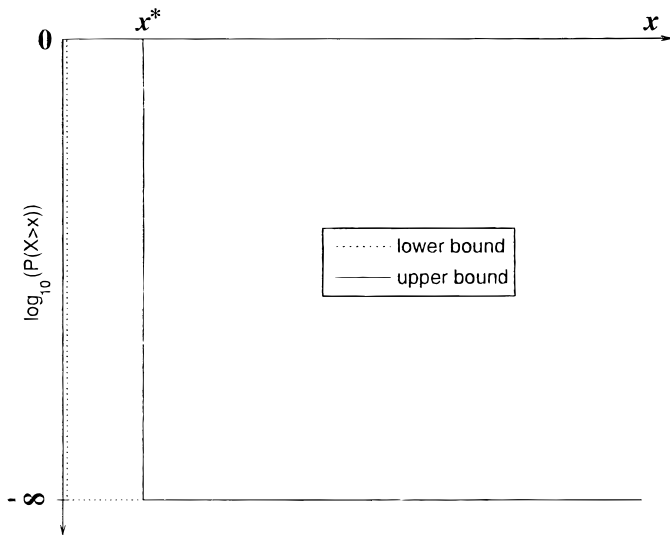


FIGURE 4. Logarithm of the bounds as a function of  $x$  for the limiting case.

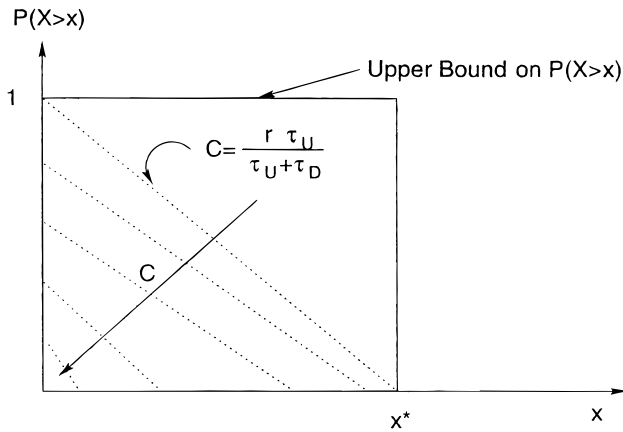


FIGURE 5. The deterministic on-off source probabilities and bounds.

capacity  $c_1$ . The output from the buffer acts as an input to another infinite-capacity buffer whose output capacity is  $c_2$ .

The effective bandwidth of the exponential on-off source is

$$eb(v) = \frac{rv - \alpha - \beta + \sqrt{(rv - \alpha - \beta)^2 + 4\beta rv}}{2v}. \tag{108}$$

Assume

$$r\beta/(\alpha + \beta) < c_2 < c_1 < r.$$

We study the buffer-content processes of the respective buffers  $\{X_1(t), t \geq 0\}$  and  $\{X_2(t), t \geq 0\}$ . See Figure 6 for an illustration of the model.

For the exponential on-off source we have

$$P\{X_1 > x\} = \frac{r\beta}{c_1(\alpha + \beta)} e^{-\eta_1 x}, \tag{109}$$

where

$$\eta_1 = \frac{c_1 \alpha + c_1 \beta - \beta r}{c_1(r - c_1)}.$$

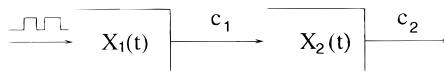


FIGURE 6. Exponential on-off input to buffers in tandem.



Using the analysis in Narayanan [23], we model the output process from the first buffer as an alternating renewal process. Thus, the input source to the second buffer can be modeled as a general on-off source with an on-time distribution  $U(t)$  (with mean  $r/(c_1(\alpha + \beta) - r\beta)$ ), off-time distribution  $D(t)$  (with mean  $1/\beta$ ), and rate  $c_1$ , where

$$U(t) = \left( \frac{a_2^2 a_3}{2a_1} \right) \sum_{k=0}^{\infty} \left[ \left( \frac{a_2}{2a_1} \right)^{2k} \frac{(2k)!}{k!(k+1)!} \right] \left[ 1 - \sum_{n=0}^{2k} \left( \frac{e^{-a_1 t/r} (a_1 t/r)^n}{n!} \right) \right], \tag{110}$$

with  $a_1 = \beta r - \beta c_1 + \alpha c_1$ ,  $a_2 = \sqrt{4\alpha\beta c_1(r - c_1)}$ ,  $a_3 = 1/(2\beta(r - c_1))$ , and

$$D(t) = 1 - e^{-\beta t}.$$

The LST of the distribution  $U(\cdot)$  is

$$\tilde{U}(w) = \begin{cases} \frac{w + \beta + c_1 s_0(w)}{\beta} & \text{if } w \geq w^* \\ \infty & \text{otherwise,} \end{cases}$$

where  $w^* = (2\sqrt{c_1\alpha\beta(r - c_1)} - r\beta - c_1\alpha - c_1\beta)/r$ ,  $s_0(w) = (-b - \sqrt{b^2 + 4w(w + \alpha + \beta)c_1(r - c_1)})/(2c_1(r - c_1))$  and  $b = (r - 2c_1)w + (r - c_1)\beta - c_1\alpha$ . The LST of the distribution  $D(\cdot)$  is

$$\tilde{D}(w) = \begin{cases} \frac{\beta}{\beta + w} & \text{if } w > -\beta \\ \infty & \text{otherwise.} \end{cases}$$

From Kulkarni and Gautam [19] we have the effective bandwidth of this general on-off source,  $eb_2(v)$ , given by

$$eb_2(v) = \begin{cases} eb_1(v) & \text{if } 0 \leq v \leq v^* \\ (eb_1(v^*) - c_1) \frac{v^*}{v} + c_1 & \text{if } v > v^*, \end{cases}$$

where

$$v^* = \frac{\beta}{r} \left( \sqrt{\frac{c_1\alpha}{\beta(r - c_1)}} - 1 \right) + \frac{\alpha}{r} \left( 1 - \sqrt{\frac{\beta(r - c_1)}{c_1\alpha}} \right) \tag{111}$$

and  $eb_1(v)$  is from Eq. (108). Note that  $\eta_2$  is obtained by solving

$$eb_2(\eta_2) = c_2.$$

If  $\eta_2 \leq v^*$ , we have,

$$\Phi(\eta_2) = \begin{bmatrix} 0 & \tilde{D}(\eta_2 c_2) \\ \tilde{U}(-\eta_2(c_1 - c_2)) & 0 \end{bmatrix},$$

and  $e(\Phi(\eta_2)) = 1$  (note that if  $\eta_2 > v^*$ , then  $e(\Phi(\eta_2)) = 1$  has no solutions). Then we obtain  $h_2$  by solving

$$[1 \ h_2]\Phi(\eta_2) = [1 \ h_2]$$

as  $h_2 = \tilde{D}(\eta_2 c_2)$ .

Intuitively, a random variable with the distribution  $U(t)$  (of Eq. (110)) is a Decreasing Failure Rate random variable (since  $U(t)$  represents the busy period distribution). The intuition can be verified (after a lot of algebra) using the expression for  $U(t)$  in Eq. (110).

Using Theorem 1, the steady-state distribution of the buffer-content process  $\{X_2(t), t \geq 0\}$  is bounded as

$$C_{2*} e^{-\eta_2 x} \leq P(X_2 > x) \leq C_2^* e^{-\eta_2 x},$$

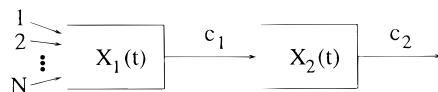
where

$$C_2^* = \frac{\frac{\tilde{D}(\eta_2 c_2) - 1}{-\eta_2 c_2} + \frac{\tilde{U}(-\eta_2(c_1 - c_2)) - 1}{\eta_2(c_1 - c_2)} h_2}{\frac{h_2 c_1(\alpha + \beta)}{\beta(c_1(\alpha + \beta) - r\beta)} \lim_{x \rightarrow \infty} \left\{ \frac{\int_x^\infty e^{\eta_2(c_1 - c_2)y} dU(y)}{\int_x^\infty e^{\eta_2(c_1 - c_2)x} dU(y)} \right\}}, \tag{112}$$

$$C_{2*} = \frac{\frac{\tilde{D}(\eta_2 c_2) - 1}{-\eta_2 c_2} + \frac{\tilde{U}(-\eta_2(c_1 - c_2)) - 1}{\eta_2(c_1 - c_2)} h_2}{\frac{h_2 c_1(\alpha + \beta)}{\beta(c_1(\alpha + \beta) - r\beta)} \tilde{U}(-\eta_2(c_1 - c_2))}. \tag{113}$$

### 7.4. Tandem Buffers—Markov Modulated On-Off Sources

Consider the tandem buffers model in Figure 7. Input to the first buffer is from  $N$  independent and identical exponential on-off sources with on-time parameter  $\alpha$ , off-time parameter  $\beta$ , and rate  $r$ . The output from buffer 1 is directly fed into buffer 2. The output capacities of buffer 1 and 2 are  $c_1$  and  $c_2$ , respectively. We study the limiting distributions of the contents of the two buffers.



**FIGURE 7.** Tandem buffers model with multiple sources.

*Buffer 1:* Let  $Z_1(t)$  be the number of sources that are in the on state at time  $t$ . Clearly,  $\{Z_1(t), t \geq 0\}$  is an SMP (more specifically, a CTMC). Assume

$$\frac{Nr\beta}{\alpha + \beta} < c_1 < Nr.$$

We can show that (Gautam [13])  $\Phi(\delta)$  is given by

$$\phi_{ij}(\delta) = \begin{cases} \frac{i\alpha}{i\alpha + (N - i)\beta - (ir - c_1)\delta} & \text{if } j = i - 1 \\ \frac{(N - i)\beta}{i\alpha + (N - i)\beta - (ir - c_1)\delta} & \text{if } j = i + 1 \\ 0 & \text{otherwise,} \end{cases}$$

and  $e(\Phi(\delta)) = 1$  always has solutions. Using the expression for  $eb(v)$  in Eq. (108) and solving for  $\eta_1$  in  $Neb(\eta_1) = c_1$  we get

$$\eta_1 = \frac{N(c_1\alpha + c_1\beta - N\beta r)}{c_1(Nr - c_1)}.$$

The eigenvectors are obtained by solving

$$h = h\Phi(\eta_1).$$

The limiting distribution of the buffer-content process  $\{X_1(t) | t \geq 0\}$  is given by

$$C_{1*}e^{-\eta_1 x} \leq P\{X_1 > x\} \leq C_1^*e^{-\eta_1 x},$$

where

$$C_1^* = \frac{\sum_{i=0}^N \frac{h_i}{\eta_1(ir - c_1)} \left( \sum_{j=0}^N (\phi_{ij}(\eta_1)) - 1 \right)}{\min_{i:ir > c_1} \frac{h_i}{p_i} \frac{1}{i\alpha + (N - i)\beta - \eta_1(ir - c_1)}},$$

$$C_{1*} = \frac{\sum_{i=0}^N \frac{h_i}{\eta_1(ir - c_1)} \left( \sum_{j=0}^N (\phi_{ij}(\eta_1)) - 1 \right)}{\max_{i:ir > c_1} \frac{h_i}{p_i} \frac{1}{i\alpha + (N - i)\beta - \eta_1(ir - c_1)}},$$

and

$$p_i = \frac{a_i \tau_i}{\sum_{m=0}^N a_m \tau_m} = \frac{N!}{i!(N - i)!} \frac{\alpha^{N-i} \beta^i}{(\alpha + \beta)^N}.$$

Buffer 2: Let  $M = \lceil c_1/r \rceil$ . Define

$$Z_2(t) = \begin{cases} Z_1(t) & \text{if } X_1(t) = 0 \\ M & \text{if } X_1(t) > 0, \end{cases} \tag{114}$$

where  $Z_1(t)$  is the number of sources on at time  $t$ . Let  $R_1(t)$  be the output rate from the first buffer at time  $t$ . We assume that

$$\frac{Nr\beta}{\alpha + \beta} < c_2 < c_1.$$

We can see that the  $\{Z_2(t), t \geq 0\}$  process is an SMP on state space  $\{0, 1, \dots, M\}$  with kernel

$$G(t) = [G_{ij}(t)]$$

derived below. For  $i = 0, 1, \dots, M - 1$  and  $j = 0, 1, \dots, M$ , let

$$G_{ij}(t) = \begin{cases} \frac{i\alpha}{i\alpha + (N - i)\beta} (1 - \exp\{-(i\alpha + (N - i)\beta)t\}) & \text{if } j = i - 1 \\ \frac{(N - i)\beta}{i\alpha + (N - i)\beta} (1 - \exp\{-(i\alpha + (N - i)\beta)t\}) & \text{if } j = i + 1 \\ 0 & \text{otherwise.} \end{cases}$$

To describe  $G_{Mj}(t)$ , we need to define the first passage time in  $\{X_1(t), t \geq 0\}$  process as follows:

$$T = \min\{t > 0 : X_1(t) = 0\}.$$

Then, for  $j = 0, 1, \dots, M - 1$ , we have

$$G_{Mj}(t) = P\{T \leq t, Z_1(T) = j | X_1(0) = 0, Z_1(0) = M\}.$$

Note that  $G_{MM}(t) = 0$ . The Laplace–Stieltjes transform (LST) of  $G_{Mj}(t)$  can be computed using the analysis in Narayanan and Kulkarni [24]. See Gautam [13] for a detailed derivation.

From Kulkarni and Gautam [19] we have the effective bandwidth of the output from Buffer 1,  $eb_2(v)$ , given by

$$eb_2(v) = \begin{cases} N \, eb_1(v) & \text{if } 0 \leq v \leq v^* \\ (N \, eb_1(v^*) - c) \frac{v^*}{v} + c & \text{if } v > v^*, \end{cases}$$

where  $eb_1(v)$  is from Eq. (108) and

$$v^* = \frac{\beta}{r} \left( \sqrt{\frac{c_1 \alpha}{\beta(Nr - c_1)}} - 1 \right) + \frac{\alpha}{r} \left( 1 - \sqrt{\frac{\beta(Nr - c_1)}{c_1 \alpha}} \right). \tag{115}$$

Hence, solving

$$eb_2(\eta_2) = c_2,$$

we get

$$\eta_2 = \min \left\{ \frac{N(c_2\alpha + c_2\beta - N\beta r)}{c_2(Nr - c_2)}, \frac{h(v^*) - c_1v^*}{c_2 - c_1} \right\},$$

where

$$h(v^*) = \frac{(rv^* - \alpha - \beta + \sqrt{(rv^* - \alpha - \beta)^2 + 4\beta rv^*})N}{2}. \tag{116}$$

If  $\eta_2 \leq v^*$ , then

$$\phi_{ij}(\eta_2) = \begin{cases} \tilde{G}_{ij}(-\eta_2(ir - c_2)) & \text{if } 0 \leq i \leq M - 1, \\ \tilde{G}_{ij}(-\eta_2(c_1 - c_2)) & \text{if } i = M, \end{cases}$$

and  $e(\Phi(\eta_2)) = 1$ . We obtain  $h$  by solving

$$h\Phi(\eta_2) = h.$$

It can be shown that the random variables associated with the distribution  $G_{Mj}(x)/G_{Mj}(\infty)$  have a decreasing failure rate. Hence, from Theorem 7,  $\Psi_{\min}(M, j)$  and  $\Psi_{\max}(M, j)$  occur at  $x = \infty$  and  $x = 0$ , respectively. Thus, using Theorem 1, we can find bounds for the steady-state distribution of the buffer-content process  $\{X_2(t), t \geq 0\}$ .

In Figure 8 we illustrate the upper and lower bounds on the limiting distribution of the buffer-content process

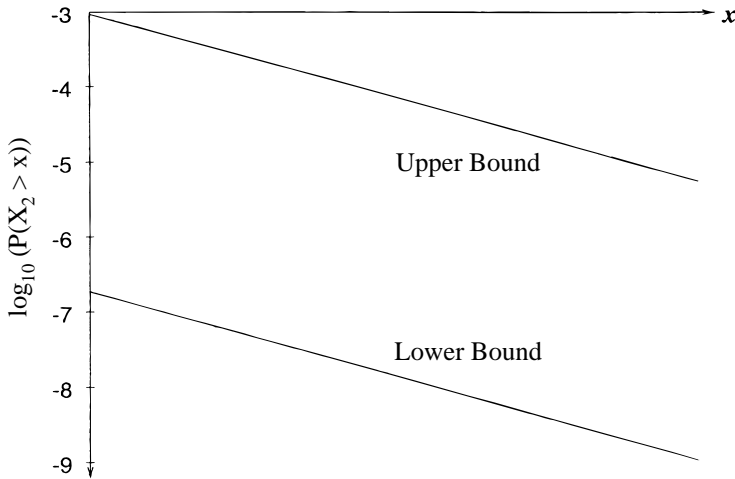
$$\lim_{t \rightarrow \infty} P\{X_2(t) > x\} = P\{X_2 > x\}$$

for a numerical example with  $\alpha = 1, \beta = 0.3, r = 1, c_1 = 13.22, c_2 = 10.71$ , and  $N = 16$ .

In a forthcoming paper, we will discuss the case  $\eta_2 > v^*$  (when  $e(\Phi(v)) = 1$  has no solutions) and use the results to solve Quality of Service problems in multipriority networks.

### 8. CONCLUSIONS AND EXTENSIONS

In this paper we have derived upper and lower bounds on the steady-state distribution of the buffer-content process (with constant output capacity) whose input is modulated by a single SMP. Unlike the effective-bandwidth approximation, where the results are valid only in the asymptotic case of large buffers and small tail probabilities, the SMP bounds are valid for all buffer sizes and tail probabilities. We have also illustrated how to compute the SMP bounds and obtain closed form algebraic



**FIGURE 8.** The upper and lower bounds as a function of  $x$ .

expressions for certain special cases, as well as reported the results (upper and lower bounds) using several examples.

We have derived the bounds for the buffer-content process when the input to the buffer is generated from several sources, each modulated by an independent SMP. Since the result is very similar to the single SMP source, we do not go into the details of the proof of the theorem. We conclude, however, that the only significant difference is the minimization or the maximization being done over a slightly different set. We consider examples and applications of Theorem 6 in detail in a few other papers dealing with multiplexing multiclass traffic in high-speed telecommunication networks.

It may be possible to extend the SMP bounds result to more general sources like MRGP sources and general Markovian sources, and with time-varying output capacities.

We also believe that  $P(X > x) \sim Ce^{-\eta x}$  for some constant  $C > 0$  (perhaps under the mild conditions); however, this problem for more than one stream seems to be very difficult. In fact, we think that  $W(\tau(x)) \xrightarrow{D} W^0$  for some variable  $W^0$ . Then, from Eq. (96) and the fact that function  $v$  is bounded and continuous, we would obtain the needed assertion.

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